

# A Structural Method for Hamiltonian Graphs

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Abstract

In this paper, we shall introduce a special structure for graphs and show that a graph  $G$  is hamiltonian if and only if  $G$  has such a special structure. Using this result, we can prove a new weakened version of Fan's condition for hamiltonian graphs, which generalizes a recent result of Bedrossian, Chen and Schelp (1993).

## 1 Preliminaries and Main Results

We consider only finite undirected graphs without loops or multiple edges. The set of vertices of a graph  $G$  is denoted by  $V(G)$  or just by  $V$ ; the set of edges by  $E(G)$  or just by  $E$ . We use  $|G|$  as a symbol for the cardinality of  $V(G)$ . If  $H$  and  $S$  are subsets of  $V(G)$  or subgraphs of  $G$ , we denote by  $N_H(S)$  the set of vertices in  $H$  which are adjacent to some vertex in  $S$ , and set  $d_H(S) = |N_H(S)|$ . If  $S = \{u\}$  and  $H = G$ , then let  $N_G(u) = N(u)$  and set  $d_G(u) = d(u)$ . For  $D \subseteq V(G)$ ,  $G[D]$  denotes the subgraph of  $G$  induced by  $D$ . For basic graph-theoretic terminology, we refer the reader to [3].

**Definition 1.** Let  $H$  be a subgraph of  $G$  and  $x, y \in V(G) \setminus V(H)$ .  $\{x, y\}$  is called a pair of useful vertices of  $H$  if  $G[H \cup \{x, y\}]$  contains a hamiltonian path connecting  $x$  and  $y$ .

**Definition 2.** A graph  $G$  is called  $L$ -decomposable if  $G$  can be separated into  $k + 1$  pairwise disjoint subgraphs  $G_0, G_1, \dots, G_k$  such that the following four conditions are satisfied:

- 1)  $G_0$  is complete.
- 2) For any  $1 \leq i \leq k$ , there exists a subset  $S_i \subseteq N_{G_0}(G_i)$  with at least two vertices which contains a vertex  $x$  such that for every  $y (\neq x) \in S_i$ ,  $\{x, y\}$  is a pair of useful vertices of  $G_i$ .
- 3) For any three distinct  $S_i, S_j, S_l$ , we have  $S_i \cap S_j \cap S_l = \emptyset$ .

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4) For any positive integer  $r \leq k$ ,  $|\bigcup_{1 \leq j \leq r} S_j| = r$  if and only if  $|V(G_0)| = k = r$ .

If  $G$  is L-decomposable, then we say the partition  $G_0, G_1, \dots, G_k$  which satisfies the four conditions above a *L-decomposition of  $G$* . In Section 2, we shall prove the following structural theorem.

**Theorem 1.** A graph  $G$  is hamiltonian if and only if  $G$  has a L-decomposition.

Theorem 1 has some applications. We shall give some examples here. In order to do this, we need some additional terminology and notations.

In Figure 1, we define four kinds of graphs, C-graph, F-graph, B-graph and N-graph.

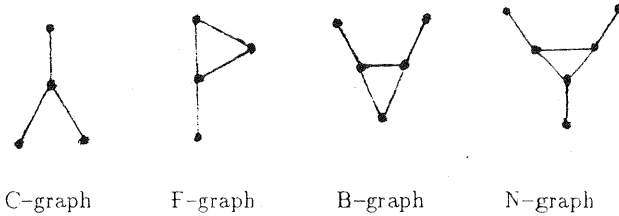


Figure 1.

Let  $S, T$  be two induced subgraphs of  $G$  with  $\max\{|S|, |T|\} < |G|$ . A graph  $G$  of order  $n$  is said to satisfy property  $ST(n)$  if for any pair of vertices  $x$  and  $y$  at distance two in  $S$  or  $T$ ,  $\max\{d(x), d(y)\} \geq n/2$ . If  $G$  contains no  $S$  as an induced subgraph, we call  $G$   $S$ -free. If  $G$  contains neither  $S$  nor  $T$  as an induced subgraph, we call  $G$   $ST$ -free.

The closure of a graph  $G$  denoted by  $\bar{G}$ , is the graph obtained from  $G$  by recursively joining pairs of nonadjacent vertices whose degree sum is at least  $|V(G)|$  until no such pair remains. Let  $V_0 = \{x : d(x) \geq n/2, x \in V(G)\}$ .

The following result is due to Bondy and Chvátal.

**Theorem 2[2].** A graph  $G$  is hamiltonian if and only if  $\bar{G}$  is hamiltonian.

Now, using Theorems 1 and 2, we can easily prove the following two theorem known before.

**Theorem 3[4].** Let  $G$  be a 2-connected graph of order  $n$ . If each pair of vertices  $x$  and  $y$  at distance 2 satisfies  $\max\{d(x), d(y)\} \geq n/2$ , then  $G$  is hamiltonian.

**Theorem 4[1].** Let  $G$  be a 2-connected graph of order  $n$ . If  $G$  satisfies property  $CF(n)$ , then  $G$  is hamiltonian.

To prove Theorems 3 and 4, we assume, by contradiction, that  $G$  is a counterexample with as many as possible edges. By Theorem 2,  $G[V_0]$  is a complete subgraph of  $G$ . Let  $G_0$  be an induced complete subgraph of  $G$  with as many as possible vertices and  $V_0 \subseteq V(G_0)$ . Let  $G_1, G_2, \dots, G_k$  be the components of  $G \setminus G_0$ . We can easily

verify that  $G_0, G_1, \dots, G_k$  is a L-decomposition of  $G$  under conditions of Theorem 3 or Theorem 4, which leads to a contradiction by Theorem 1.

In section 3, we shall prove the following more general theorem by using Theorems 1 and 2.

**Theorem 5.** Let  $G$  be a 2-connected graph of order  $n$ . If  $G$  satisfies property  $CB(n)$ , then  $G$  is hamiltonian.

## 2 The Proof of Theorem 1

If  $G$  is a hamiltonian graph, let  $C = c_1 c_2 \dots c_n c_1$  be a hamiltonian cycle of  $G$ . Set  $G_0 = G[\{c_1, c_2\}]$  and  $G_1 = G[\{c_3, \dots, c_n\}]$ . Then  $G_0, G_1$  satisfy the four conditions of Definition 2. Thus  $G$  has a L-decomposition.

Conversely, let  $G_0, G_1, \dots, G_k$  be a L-decomposition of  $G$ . By Definition 2,  $G_0$  is a complete subgraph with  $|G_0| \geq 2$  and for any  $1 \leq i \leq k$ , there exists some  $S_i \subseteq N_{G_0}(G_i)$  which satisfies the conditions 2)–4) of Definition 2. By condition 2),  $S_i$  contains a vertex  $x_i$  such that for any  $y \in S_i \setminus \{x_i\}$ ,  $\{x_i, y\}$  is a pair of useful vertices of  $G_i$  for all  $1 \leq i \leq k$ . Using the following Claim we will give a structural proof of the sufficiency.

**Claim.**  $G_0$  contains either a cycle  $C = u_1 u_2 \dots u_k u_1$  with  $|V(C)| = |G_0|$  (when  $|G_0| = 2$ ,  $C$  is just an edge.) such that

$$\{u_i, u_{j+1}\} = \{x_i, y_j\}, \quad j = 1, \dots, k, \quad j \bmod k \quad (*)$$

or  $q$  pairwise disjoint paths  $P_i = u_1 u_2 \dots u_{r_i+1}$ ,  $i = 1, 2, \dots, q$

$$\{u_j, u_{j+1}\} = \{x_i, y_j\}, \quad j = 1, 2, \dots, r_i, \quad (**)$$

and

$$u_1, \dots, u_{r_i+1} \notin \bigcup_{j \notin \{1, \dots, r_i\}} S_j \quad (***)$$

where  $y_j \in S_i \setminus \{x_i\}$ .

In fact, let  $P = u_1 \dots u_{r_i+1}$  be a longest path satisfying the equation (\*\*). Then  $u_2, \dots, u_{r_i} \notin \bigcup_{j \notin \{1, \dots, r_i\}} S_j$  by condition 3). If  $u_1, u_{r_i+1} \notin \bigcup_{j \notin \{1, \dots, r_i\}} S_j$ , then  $P$  is desired. Otherwise, there exists a subset, say  $S_{i, r_i+1}$ , such that  $\{u_1, u_{r_i+1}\} \cap S_{i, r_i+1} \neq \emptyset$ . By the maximality of  $P$  and  $|S_{i, r_i+1}| \geq 2$ , we have that  $S_{i, r_i+1} = \{u_1, u_{r_i+1}\}$ . Since  $|\bigcup_{1 \leq j \leq r} S_j| \geq r$  for any  $r \leq k$ , we need only to consider the following two cases.

**Case 1.**  $|\bigcup_{1 \leq j \leq r_i+1} S_j| = r_i + 1$ .

Then  $|V(G_0)| = k = r_i + 1$  by condition 4). Thus  $C = u_1 \dots u_{r_i+1} u_1$  is a cycle of  $G_0$  with  $|V(C)| = |G_0|$  satisfying (\*).

**Case 2.**  $|\bigcup_{1 \leq j \leq r_i+1} S_j| > r_i + 1$ .

By condition 3), there is a  $l \in \{1, \dots, r_i\}$  such that  $|S_l| \geq 3$ . We assume without loss of generality that  $\{x_l, y_l, z_l\} \subseteq S_l$  satisfying  $x_l = u_l$ ,  $y_l = u_{l+1}$  and  $z_l \notin$

$V(P)$ . Then we can construct a new path  $P' = z_i u_i u_{i-1} \dots u_i u_{i+r+1} u_r \dots u_{i+1}$  which is longer than  $P$  and satisfies (\*\*) when the subscripts are rewritten. This contradiction completes the proof of the Claim.

Now, from the Claim above, if  $G_0$  contains a cycle  $C$  with  $V(C) = |G_0|$  satisfying (\*), then it is easy to check that  $G$  is hamiltonian. Otherwise, by the Claim above,  $G_0$  contains  $q$  pairwise disjoint paths  $P_i = u_{i_1} u_{i_2} \dots u_{i_{r_i+1}}$ ,  $i = 1, 2, \dots, q$  which satisfy both (\*\*) and (\*\*\*), and we have  $\sum_{i=1}^q r_i = k$ . Since  $G_0$  is a complete subgraph of  $G$ , we can easily check that  $G$  has a hamiltonian cycle.

Therefore, Theorem 1 is true.  $\diamond$

Theorem 1 has the following consequence.

**Corollary 1.** Let  $G_0$  be a complete subgraph of  $G$  with  $|G_0| \geq 2$ . If  $G_0$  contains a pair of useful vertices of each component of  $G \setminus G_0$  and  $G[N(G_0)]$  is C-free, then  $G$  is hamiltonian.

**Proof.** Let  $G_1, \dots, G_k$  be all the components of  $G \setminus G_0$  and set  $G^* = G[N(G_0)]$ . By Theorem 1, it is sufficient to show that  $G_0, G_1, \dots, G_k$  is a L-decomposition of  $G$ .

By the hypothesis, we can choose  $S_i \subseteq V_{G_0}(G_i)$  such that  $S_i$  satisfies 2) of Definition 2 and  $|S_i|$  is as large as possible. Since  $G^*$  is C-free, 3) of Definition 2 is satisfied. Thus we only need to show that 4) of Definition 2 is also satisfied.

In fact, let  $r \leq k$  be any positive integer. Since  $G^*$  is C-free, we have  $|V(G_0)| \geq k \geq r$ . If  $|V(G_0)| = k = r$ , then  $|\bigcup_{1 \leq j \leq r} S_{i_j}| = r$ . Conversely, if  $|\bigcup_{1 \leq j \leq r} S_{i_j}| = r$ , then  $|S_{i_j}| = 2$  ( $j = 1, 2, \dots, r$ ) and each vertex  $x \in \bigcup_{1 \leq j \leq r} S_{i_j}$  is a common vertex of some two pairs of useful vertices. Let  $x \in S_{i_r} \cap S_{i_t}$  and  $y \in N_{G_{i_r}}(x)$ ,  $z \in N_{G_{i_t}}(x)$ . When  $|G_0| > r$ , then there exists some  $w \in V(G_0) \setminus (\bigcup_{1 \leq j \leq r} S_{i_j})$ . Since  $G^*$  is C-free, we have  $wy \in E$  or  $wz \in E$ . Therefore, either  $S_{i_r} \cup \{w\}$  or  $S_{i_t} \cup \{w\}$  still satisfies 2) of Definition 2, which contrary to the choice of  $S_{i_r}$  or  $S_{i_t}$ . Thus  $|V(G_0)| = k = r$ . This completes the proof of Corollary 1.  $\diamond$

### 3 The Proof of Theorem 5

In order to prove Theorem 5, we need the following theorem.

**Theorem 6[5].** If  $G$  is 3-connected and CN-free, then for any distinct vertices  $x, y$  of  $G$ , there exists a hamiltonian path connecting  $x$  and  $y$ .

Now, set  $V_0 = \{x \in V(G) : d(x) \geq n/2\}$ . By Theorem 2, we may assume that  $G[V_0]$  is a complete subgraph of  $G$  if  $V_0 \neq \emptyset$ . Let  $G_0$  be a complete subgraph of  $G$  such that  $V_0 \subseteq V(G_0)$  and  $|V(G_0)|$  is as large as possible. Let  $G_1, \dots, G_k$  be all the components of  $G \setminus G_0$ . Then by the property  $CB(n)$ ,  $G[N(G_0)]$  is C-free and  $G_s$  is CB-free for any  $1 \leq s \leq k$ . By Corollary 1, we need only to show that  $G_0$  contains a pair of useful vertices of  $G_s$  for  $1 \leq s \leq k$ .

Assume that there is a component  $G_s$  of  $G \setminus G_0$  such that  $G_0$  does not contain

any pair of useful vertices of  $G_s$ . Let  $S$  be a minimal cut vertex set of  $G_s$  and  $v \in S$ . Then by the assumption and Theorem 6,  $|S| \leq 2$ . Since  $G_s$  is C-free,  $G_s \setminus S$  has only two components  $H_1, H_2$ . Let  $H = G[V(H_1) \cup V(H_2) \cup \{v\}]$  and  $S_{-i} = \{u \in V(H_1) : d_H(u, v) = i\}$  and  $S_i = \{u \in V(H_2) : d_H(u, v) = i\}$  for  $i \geq 0$ . Denote  $m := \max\{i : S_i \neq \emptyset\}$  and  $n := \max\{i : S_{-i} \neq \emptyset\}$ . Clearly, we have  $V(G_s) = S \cup (\bigcup_{i=-n}^m S_i)$ , and  $G[S_i \cup S_j]$  is complete if and only if  $|i - j| = 1$  since  $G_s$  is CB-free.

If  $|S| = 1$ , then there exist some  $x \in S_m$  and  $y \in S_{-n}$  such that neither  $x$  nor  $y$  is a cut vertex of  $G_s$  and  $N_{G_0}(x) \neq \emptyset$  and  $N_{G_0}(y) \neq \emptyset$ , since  $G$  is 2-connected. Because of the structure of  $G_s$ , there exists a path  $P$  connecting  $x$  and  $y$  in  $G_s$  with  $V(P) = V(G_s)$ . Thus by the assumption,  $N_{G_0}(x) = N_{G_0}(y)$  and  $|N_{G_0}(x)| = 1$ , which is contrary to the fact that  $G[N(G_0)]$  is C-free.

If  $|S| = 2$ , let  $v' \in S$  and  $v' \neq v$ . Since  $G_s$  is 2-connected,  $N(v') \cap S_i \neq \emptyset$  for some  $1 \leq i \leq m$  and  $N(v') \cap S_{-j} \neq \emptyset$  for some  $1 \leq j \leq n$ . Let  $i_0 = \max\{i : N(v') \cap S_i \neq \emptyset\}$  and  $j_0 = \max\{j : N(v') \cap S_{-j} \neq \emptyset\}$ . By the hypothesis of Theorem 5, we may assume that there exists some  $t$  with  $0 \leq t \leq m$  such that  $N_{G_0}(S_t) \neq \emptyset$ .

Since  $G_s$  is 2-connected, we have

(a)  $|S_i| \geq 2$  for any  $m - 1 \geq i \geq i_0$  and  $|S_{-j}| \geq 2$  for any  $n - 1 \geq j \geq j_0$ .

By (a) and the structure of  $G_s$ , we have

(b) If  $|S_m| \geq 2$ , then for any two distinct vertices  $x$  and  $y$  in  $S_m$ , there exists a path  $P$  in  $G_s$  connecting  $x$  and  $y$  with  $V(P) = V(G_s)$ .

(c) For any  $x \in S_{i-1}$  and  $y \in S_i$  ( $1 \leq i \leq m$ ), there exists a path in  $G_s$  connecting  $x$  and  $y$  with  $V(P) = V(G_s)$ .

Since  $|N_{G_0}(G_s)| \geq 2$ . By the assumption, (c) and the hypothesis of Theorem 5, we have

(d)  $n + m \geq 3$ .

Now, we distinguish the following two cases.

Case 1.  $0 \leq t < m$ , that is there exists some  $x \in S_t$  and  $y \in V(G_0)$  such that  $xy \in E$ .

Then by the hypothesis of Theorem 5 and  $1 \leq t < m$ , there exists a vertex  $z \in S_{t-1}$  or  $z \in S_{t+1}$  such that  $yz \in E$ . By the assumption and (c), for any  $y' \in V(G_0) \setminus \{y\}$  and  $w \in S_{t-1} \cup S_{t+1}$ ,  $y'w \notin E$ . Thus we can find a vertex set  $F = \{x, y, z, y', w\}$  such that  $G[F]$  is a B-graph and does not satisfy the condition of Theorem 5, a contradiction.

Case 2. For any  $0 \leq i \leq m - 1$ ,  $N_{G_0}(S_i) = \emptyset$ , that is  $t = m$ .

Symmetrically, we may assume that for any  $0 \leq j \leq n - 1$ ,  $N_{G_0}(S_{-j}) = \emptyset$ .

If  $N_{G_0}(v') \neq \emptyset$ , let  $y \in V(G_0)$  such that  $v'y \in E$ . Then by the hypothesis of Theorem 5, we have  $y \in N_{G_0}(S_{i_0})$  or  $y \in N_{G_0}(S_{-j_0})$ . Thus  $i_0 = m$  or  $j_0 = n$ .

Without loss of generality, let  $y \in N_{G_0}(S_{i_0})$ . When  $y \notin N_{G_0}(S_{-j_0})$ , then by the hypothesis of Theorem 5, there exists a vertex  $y' \in V(G_0) \setminus \{y\}$  such that  $v'y' \in E$  or  $y' \in N_{G_0}(S_m)$  or  $y' \in N_{G_0}(S_{-j_0})$  whenever  $j_0 = n$ . By the structure of  $G_s$ , we can derive that  $\{y, y'\}$  is a pair of useful vertices of  $G_s$ , contrary to the assumption. When  $y \in N_{G_0}(S_{i_0}) \cap N_{G_0}(S_{-j_0})$ , that is  $i_0 = m$  and  $j_0 = n$ . Since  $G$  is 2-connected, there exists a vertex  $y' \in V(G_0)$  such that  $y' \in N_{G_0}(S_m) \cup N_{G_0}(S_{-n})$  or  $v'y' \in E$ . Also by the structure of  $G_s$ , we can derive that  $\{y, y'\}$  is a pair of useful vertices of  $G_s$ , contrary to the assumption. Hence in rest proof we suppose that  $N_{G_0}(v') = \emptyset$ .

Since  $G$  is 2-connected, there exist  $x \neq x' \in S_m \cup S_{-n}$  and  $y \neq y' \in V(G_0)$  such that  $xy \in E$  and  $x'y' \in E$ . By the assumption and (b),  $\{x, x'\} \not\subseteq S_m$  and  $\{x, x'\} \not\subseteq S_{-n}$ . Let  $x \in S_m$  and  $x' \in S_{-n}$ . By (d), let  $m \geq 2$ , then  $S_m \in N(y)$  by the hypothesis of Theorem 5.

If  $i_0 = m$ , then  $S_{m-1} \subseteq N(v')$  by the hypothesis of Theorem 5. Thus by the structure of  $G_s$ , we can derive that  $\{y, y'\}$  is a pair of useful vertices of  $G_s$ , contrary to the assumption. If  $i_0 < m$ , then  $S_{i_0-1} \subseteq N(v')$ . Thus we can also get a contradiction as before.

Therefore, Theorem 5 is true.

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