

by

MARTA SVED and ROBERT J. CLARKE
Department of Pure Mathematics
The University of Adelaide

Abstract: The infinite chessboard, or equivalently the lattice of points of integer coordinates in the plane, gives rise to paths consisting of steps between neighbouring lattice-points. On the chessboard the king is allowed horizontal, vertical and diagonal steps, one at a time. The number of possible paths from the origin to some arbitrary lattice point has been found by various authors. In the present work the number of paths modulo a prime p is explored. Under conditions placed on the allowable steps, the arrays representing the number of paths mod p show fractal structures which indicate that once these numbers are known for the fundamental region of lattice points with coordinates between 0 and p , there are simple formulae or algorithms for determining these numbers for arbitrary lattice points. The situation is similar to that known for binomial coefficients and summarized by Lucas' theorem.

INTRODUCTION

The language devised for arrays representing binomial coefficients modulo p , of the structure illustrated in Figure 1 below, will be used.

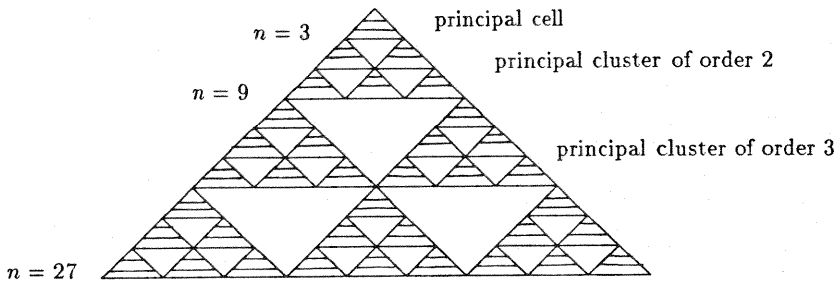


Figure 1: Binomial array $\binom{n}{k} \pmod{3}$

Detailed descriptions of these arrays can be found in [2] or [3].

The triangles shaded in the diagram are *cells*, similar to the principal cell in the sense that the entries in each are in constant ratio to the corresponding entries of the principal cell, determined by the entry at the *head* of the cell. In the same way, clusters are similar to the principal cluster of the same order.

The inverted blank triangles represent *zero holes*, which are arrays consisting of zero entries exclusively.

Each *cluster* of order m consists of p *layers*, the layers containing clusters of order $m - 1$, alternating with zero-holes.

Although the structure of the arrays to be discussed in the following differs in some respects from that of the binomial array, the terms cell, cluster, head of cell or cluster will be used, with specific notation for the entries introduced in each case.

2. KING'S WALK, PROGRESS RESTRICTED TO POSITIVE DIRECTION

In a recent paper, M. Razpet [1] deduced formulae for the possible number of walks $w(i, j)$ from the origin to a point (i, j) ($i, j \geq 0$). For $w(i, j) \pmod p$ he obtained a result analogous to that of Lucas for binomials:

Let

$$i = a_m p^m + \dots + a_1 p + a_0,$$

$$j = b_m p^m + \dots + b_1 p + b_0$$

(where $0 \leq a_s, b_s < p$ for $0 \leq s \leq m$) be the expansions to base p of i and j .

Let

$$w(i, j) \pmod p = \bar{w}(i, j).$$

Then

$$\bar{w}(i, j) = \bar{w}(a_m, b_m) \bar{w}(a_{m-1}, b_{m-1}) \dots \bar{w}(a_0, b_0) \quad (1)$$

Razpet obtained this identity by algebraic means and then produced computer outputs of $\bar{w}(i, j)$.

In this discussion we follow a somewhat different route. By establishing first the general formula for $w(i, j)$ from the recursion formula, we apply it to the principal cell: $0 \leq i, j < p$, and from this we show that a structure of cells and clusters is initiated. This implies the Lucas type relation (1). The cells and clusters in this situation will have square shapes instead of the triangular shapes of the binomial array, and the nature of the zero-holes will be different.

To fit with computer outputs, coordinate systems will be differently oriented from the usual way, as shown in Figure 2.

Since the king is allowed steps in the positive direction under the conditions imposed, the point (i, j) can be reached through the final steps indicated in the diagram. Hence

$$w(i, j) = w(i - 1, j - 1) + w(i - 1, j) + w(i, j - 1). \quad (2)$$

The general formula for $w(i, j)$ is devised from the recursion formula, by using *shift* operators defined by

$$\sigma_1 w(i, j) = w(i - 1, j - 1)$$

$$\sigma_2 w(i, j) = w(i - 1, j)$$

$$\sigma_3 w(i, j) = w(i, j - 1).$$

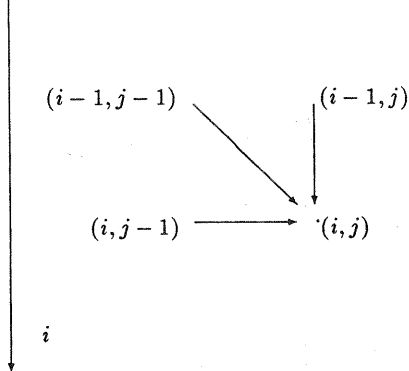


Figure 2

Thus (2) can be written in the form

$$(\sigma_1 + \sigma_2 + \sigma_3) w(i, j) = w(i, j). \quad (3)$$

We find first $w_k(i, j)$, the number of ways in which (i, j) can be reached in k steps from $(0, 0)$.

Iterating (3), we obtain that in general for k steps terminating in (i, j)

$$w(i, j) = (\sigma_1 + \sigma_2 + \sigma_3)^k w(i, j) \quad (4)$$

and

$$w(i, j) = \sum_{k=\max(i,j)}^{i+j} w_k(i, j). \quad (5)$$

Denote by d, v, h the number of diagonal, vertical and horizontal steps, respectively, making up the path; we obtain for a path consisting of k steps

$$w(i, j) = \sum \frac{k!}{d!v!h!} \sigma_1^d \sigma_2^v \sigma_3^h w(i, j) \quad (6)$$

where

$$\sigma_1^d \sigma_2^v \sigma_3^h w(i, j) = w(i - (d + v), j - (d + h)). \quad (7)$$

Since we consider the case when k steps lead from $(0, 0)$ to (i, j) , we have

$$\sigma_1^d \sigma_2^v \sigma_3^h w(i, j) = w(0, 0) = 1;$$

thus $i - (d + v) = 0$, $j - (d + h) = 0$, while $d + v + h = k$. These three equations determine precisely one value for the set (d, v, h) , namely $v = k - j$, $h = k - i$, $d = i + j - k$, provided that $\max(i, j) \leq k \leq i + j$.

From (6) and (7)

$$w_k(i, j) = \frac{k!}{(k-i)!(k-j)!(i+j-k)!} = \binom{k}{i} \binom{i}{k-j}. \quad (8)$$

This formula can also be given a combinatorial interpretation by counting the ways in which k steps from the origin to point (i, j) can be distributed into horizontal, vertical and diagonal steps.

Hence, from (5),

$$w(i, j) = \sum_{k=\max(i, j)}^{i+j} \binom{k}{i} \binom{i}{k-j}. \quad (9)$$

In particular, for $i = 0$ or $j = 0$, $w(i, j) = 1$. Also $w(i, j) = w(j, i)$, as expected.

Next we investigate the $w(i, j)$ array (mod p). Since $\bar{w}(i, j) = w(i, j) \pmod{p}$, the principal cell is the array $(\bar{w}(i, j) | 0 \leq i, j < p)$. For the last row of the cell we have, from (9),

$$\bar{w}(p-1, j) = \sum_{k=p-1}^{p-1+j} \binom{k}{p-1} \binom{p-1}{k-j} \pmod{p}$$

Since

$$\binom{k}{p-1} \equiv 0 \pmod{p} \text{ for } p \leq k < 2p-1,$$

$$\bar{w}(p-1, j) = \binom{p-1}{p-1-j} = \binom{p-1}{j} \equiv (-1)^j \pmod{p},$$

and, by symmetry,

$$\bar{w}(i, p-1) = (-1)^i. \quad (10)$$

Denote the *principal cell* by $C[0, 0]$, and, more generally, let

$$C[rp, sp] = (\bar{w}(i, j) | rp \leq i < (r+1)p, sp \leq j < (s+1)p)$$

(Here (rp, sp) represent the "head", the coordinates of the point at the top left corner.)

Theorem:

- (i) $\bar{w}(rp, sp) = \bar{w}(r, s)$
- (ii) $C[rp, sp]$ is a cell, that is

$$\bar{w}(rp+i, sp+j) = \bar{w}(r, s)\bar{w}(i, j). \quad (11)$$

Proof. We note first that if the rows and columns of the principal cell are extended to infinity, we obtain in both directions copies of the principal cell, as shown in Figure 3.

This follows immediately from the recursion formula (2) and (10); so $C[rp, 0]$ and $C[0, sp]$ are cells for all r, s .

Next assume that the structure is established down to the $(r-1)$ -th horizontal layer (hence, by symmetry, also to the $(r-1)$ -th vertical one). The induction can now proceed

1	1	1	1	1	1	1	1	1	1	1	1
1			-1	1		-1	1		-1	1	
1			1	1		1	1		1	1	
1			-1	1		-1	1		-1	1	
1	-1	1	-1	1	1	-1	1	-1	1	-1	1
1	1	1	1	1							
1			-1								
1			1								
1			-1								
1	-1	1	-1	1							
1	1	1	1	1							
1			-1								
1			1								
1			-1								
1	-1	1	-1	1							

Figure 3

along the r -th layer by assuming that for $s > r$, $C[(r-1)p, (s-1)p]$, $C[(r-1)p, sp]$ and $C[rp, (s-1)p]$ are cells with head entries $\bar{w}(r-1, s-1)$, $\bar{w}[r-1, s]$ and $\bar{w}[r, s-1]$, respectively, with corresponding last rows and columns of alternating terms; so by (2) and the induction hypothesis

$$\bar{w}(rp, sp) = \bar{w}(r-1, s-1) + \bar{w}(r-1, s) + \bar{w}(r, s-1) = \bar{w}(r, s).$$

Also the first row and first column of $C[rp, sp]$ consist of identical entries. The rest follows.

Next we consider the array $(w(i, j) | 0 \leq i, j < p^2)$, called the principal cluster of order 2, consisting of p layers of cells. We note here that the $C[(p-1)p, sp]$ and $C[rp, (p-1)p]$ cells have again alternating $+1$ and -1 entries in their last rows and columns.

We now apply an induction procedure similar to that used on cells, and proceed from clusters of order 2 to clusters of order 3 and so on, and thus establish the self similar structure of the $\bar{w}(i, j)$ array. \square

Tables (1) and (2) show these arrays modulo 3 and 7. The zero-holes on these are really cells or clusters with heads = 0.

The extension of Lucas' theorem, relation (1) as established by Razpet, follows now directly from (11). We write the expansions of i and j to base p shortly as $[a_m, a_{m-1}, \dots, a_0]$ and $[b_m, b_{m-1}, \dots, b_0]$ respectively, and see that the entry $\bar{w}(i, j)$ is inside a nest of clusters of orders $m, m-1, \dots, 1$ respectively, with heads $\bar{w}(a_m, b_m)$, $\bar{w}(a_{m-1}, b_{m-1}), \dots, \bar{w}(a_1, b_1)$. Thus, from (11), the relation (1) follows.

3. KING WALKS WITH MINIMAL NUMBER OF STEPS

The assumption now is that the king moves from the origin $(0, 0)$ to the point (n, k) , the path being covered in a minimal number of steps; if we assume without loss of generality

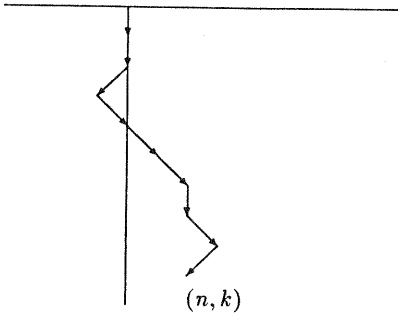


Figure 4

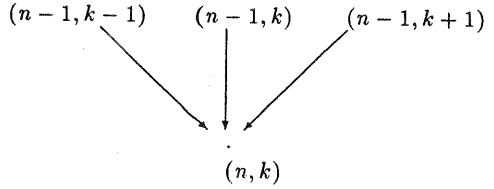


Figure 5

that $|k| \leq n$, then the minimality condition rules out horizontal steps, but diagonal steps may go in the negative direction as shown in Figure 4. Thus the last steps to reach the point (n, k) may come from one of the points shown in Figure 5.

Denote by $\left\{ \begin{matrix} n \\ k \end{matrix} \right\}$ the number of possible paths from $(0, 0)$ to (n, k) under the conditions imposed above. Figure 5 implies the recursion formula

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \left\{ \begin{matrix} n-1 \\ k-1 \end{matrix} \right\} + \left\{ \begin{matrix} n-1 \\ k \end{matrix} \right\} + \left\{ \begin{matrix} n-1 \\ k+1 \end{matrix} \right\} \quad (12)$$

with initial conditions $\left\{ \begin{matrix} 0 \\ 0 \end{matrix} \right\} = 1$ and $\left\{ \begin{matrix} 0 \\ i \end{matrix} \right\} = 0$ for $i \neq 0$.

The explicit formula for $\left\{ \begin{matrix} n \\ k \end{matrix} \right\}$ may be deduced by the shift operator method, as in (2), but for the later part it is more convenient to use the generating functions

$$f_n(x) = \sum \left\{ \begin{matrix} n \\ k \end{matrix} \right\} x^k. \quad (13)$$

The recursion (12) implies that $f_n(x) = (x + 1 + x^{-1}) f_{n-1}(x)$ and it follows from this that

$$f_n(x) = (x + 1 + x^{-1})^n. \quad (14)$$

$\left\{ \begin{matrix} n \\ k \end{matrix} \right\}$ is then the coefficient of x^k in the expansion of (14); hence

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \sum \frac{n!}{i!j!(n-i-j)!}$$

where $i - j = k$, under the restriction that $i + j \leq n$.

Let $i + j = m$, where $k \leq m \leq n$; hence $i = \frac{m+k}{2}$, $j = \frac{m-k}{2}$ and since i and j are integers, $m \equiv k \pmod{2}$. We then obtain the explicit formula

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \sum \frac{n!}{(n-m)! \left(\frac{1}{2}(n+m)\right)! \left(\frac{1}{2}(n-m)\right)!}$$

where $k \leq m \leq n$ and $m \equiv k \pmod{2}$.

An alternative formula may be obtained by setting

$$x + 1 + x^{-1} = \left(\sqrt{x} + \frac{1}{\sqrt{x}} \right)^2 - 1.$$

Then

$$\begin{aligned} f_n(x) &= \left(\left(\sqrt{x} + \frac{1}{\sqrt{x}} \right)^2 - 1 \right)^n = \sum_{r=0}^n (-1)^{n-r} \binom{n}{r} \left(\sqrt{x} + \frac{1}{\sqrt{x}} \right)^{2r} \\ &= \sum_{r=0}^n (-1)^{n-r} \binom{n}{r} \sum_{s=0}^{2r} \binom{2r}{s} x^{s-r}. \end{aligned}$$

The coefficient of x^k is then

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \sum_{r=0}^n (-1)^{n-r} \binom{n}{r} \binom{2r}{k+r}$$

An analogue for the convolution formula for binomial coefficients may be also obtained by writing, for some fixed m ,

$$(x + x^{-1} + 1)^n = (x + x^{-1} + 1)^m (x + x^{-1} + 1)^{n-m}.$$

Hence

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \sum_s \left\{ \begin{matrix} m \\ s \end{matrix} \right\} \left\{ \begin{matrix} n-m \\ k-s \end{matrix} \right\} \quad (15)$$

Tables (3) and (4) show the interesting fractal structure of $\left\{ \begin{matrix} n \\ k \end{matrix} \right\} \pmod{3}$ and $\left\{ \begin{matrix} n \\ k \end{matrix} \right\} \pmod{7}$.

In the following, the structure of $\left\{ \begin{matrix} n \\ k \end{matrix} \right\}$ will be analysed.

We begin by observing that

$$(x + 1 + x^{-1})^p \equiv x^p + 1 + x^{-p} \pmod{p}.$$

Hence

$$(x + 1 + x^{-1})^{np} \equiv (x^p + 1 + x^{-p})^n \pmod{p}. \quad (16)$$

It follows from (16) that the expansion \pmod{p} contains only those powers of x where the index is an integer multiple of p ; thus

$$\left\{ \begin{matrix} np \\ r \end{matrix} \right\} \equiv 0 \pmod{p} \quad \text{whenever } p \nmid r. \quad (17).$$

Looking at coefficients in the expansion on the right hand side of (16), we see that

$$\left\{ \begin{matrix} np \\ nk \end{matrix} \right\} \equiv \left\{ \begin{matrix} n \\ k \end{matrix} \right\} \pmod{p}. \quad (18)$$

In what follows we use again the notation

$$\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}} = \left\{ \begin{matrix} n \\ k \end{matrix} \right\} \pmod{p}.$$

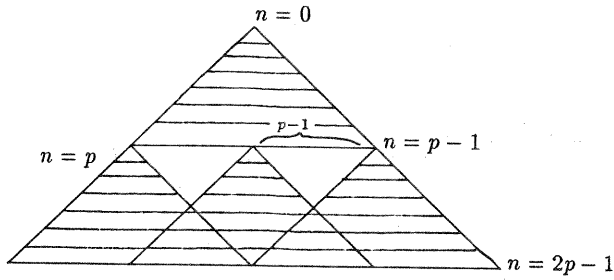


Figure 6

Next we consider the *principal cell*

$$\left(\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}} \mid 0 \leq n < p \right).$$

From the formula for $\left\{ \begin{matrix} n \\ k \end{matrix} \right\}$ it follows that

$$\left\{ \begin{matrix} n \\ n \end{matrix} \right\} = \left\{ \begin{matrix} n \\ -n \end{matrix} \right\} = 1,$$

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = 0 \text{ for } |k| > n,$$

and

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \left\{ \begin{matrix} n \\ -k \end{matrix} \right\} \text{ for all } n.$$

When $n = p$, it also follows that

$$\overline{\left\{ \begin{matrix} p \\ 0 \end{matrix} \right\}} = 1.$$

So the principal cell is a triangular structure, symmetrical about the zero axis, and having 1 for its extremal entries, and

$$1 \ 0 \ 0 \ \dots \ 1 \ 0 \ 0 \ \dots \ 0 \ 0 \ 1$$

for the entries of the row to follow (the p -th row of the $\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}}$ array).

Figure 6 shows the structure of the principal cell and the rows which follow immediately.

It follows from the recursion formula (12) that the zero entries of row p initiate zero-holes, and hence the recursion formula also implies that the first rows of the principal cells are reproduced in the shaded arrays headed by $\overline{\left\{ \begin{matrix} p \\ -p \end{matrix} \right\}}, \overline{\left\{ \begin{matrix} p \\ 0 \end{matrix} \right\}}, \overline{\left\{ \begin{matrix} p \\ p \end{matrix} \right\}} = 1$. However, the depth of the zero-holes is only $\frac{1}{2}(p-1)$, since the distance between non-zero entries decreases by one at each end with each additional row. Thus the $\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}}$ array is made more complex by the *overlap* of the cells.

On the other hand, relations (17) and (18) ensure the fractal structure of the array, with cells and zero-holes being initiated in every p -th row, while in row p^2 we have

$$\overline{\left\{ \begin{matrix} p^2 \\ kp \end{matrix} \right\}} = \overline{\left\{ \begin{matrix} p \\ k \end{matrix} \right\}} = 0 \quad \text{if } p \nmid k,$$

while

$$\overline{\left\{ \begin{matrix} p^2 \\ p^2 \end{matrix} \right\}} = \overline{\left\{ \begin{matrix} p^2 \\ -p^2 \end{matrix} \right\}} = \overline{\left\{ \begin{matrix} p^2 \\ 0 \end{matrix} \right\}} = 1;$$

so we have a structure similar to that of row p , with the two strings of zero-entries being of length $p^2 - 1$. Thus, down to $n = p^2 + \frac{p^2-1}{2}$, the principal cluster $\left(\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}} \mid n < p^2 \right)$ is reproduced in the three arrays headed by $\overline{\left\{ \begin{matrix} p^2 \\ -p^2 \end{matrix} \right\}}$, $\overline{\left\{ \begin{matrix} p^2 \\ p^2 \end{matrix} \right\}}$ and $\overline{\left\{ \begin{matrix} p^2 \\ 0 \end{matrix} \right\}}$.

As the array is developed further, clusters of higher order arise, also overlapping, so on first sight the array becomes a maze of overlapping clusters and subclusters.

In what follows, it will be shown that all $\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}}$ entries are determined by the principal cell, and while we cannot find a formula as simple as that of Lucas for binomials, and for arrays described in Section 2, an algorithm can be developed to calculate $\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}}$ in the general case.

The following theorem refers to two neighbouring cells, headed by entries in row ap , where $a > 0$. Because of the symmetry of the array and of the principal cell, we shall be looking only at the positive sides of the array and cells.

Theorem: Let the coordinates of the head of a cell be (ap, bp) , $(a, b > 0)$, and the coordinates of the right hand neighbour $(ap, (b+1)p)$, and denote the entries at the heads by ℓ, r respectively. Consider the point P having coordinates

$$n = ap + i, \quad k = bp + j \quad (0 \leq i, j < p).$$

Then

$$\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}} = \ell \overline{\left\{ \begin{matrix} i \\ j \end{matrix} \right\}} + r \overline{\left\{ \begin{matrix} i \\ p-j \end{matrix} \right\}}. \quad (19)$$

Proof: Figure 7 illustrates the region defined by the conditions imposed on i and j , with the positive side of the cell of reference, showing the entry ℓ at its head, overlapping with the negative side of its right neighbour.

We begin the proof for the "free part" of the array, when $0 \leq i \leq \frac{p-1}{2}$.

We consider three cases:

(i) $0 \leq j \leq i$

P is in the free part of the left cell; hence by similarity

$$\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}} = \ell \overline{\left\{ \begin{matrix} i \\ j \end{matrix} \right\}}$$

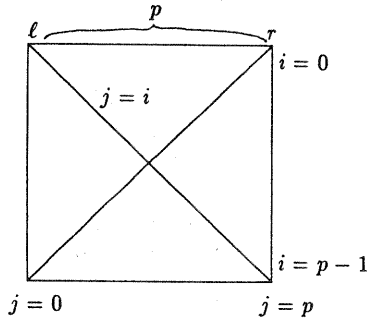


Figure 7

while $p - j > i$. Hence

$$\overline{\left\{ \begin{matrix} i \\ p-j \end{matrix} \right\}} = 0;$$

so (19) follows.

(ii) $i < j < p - i$, so $p - j > i$.

In this case

$$\overline{\left\{ \begin{matrix} i \\ j \end{matrix} \right\}} = 0 \quad \text{and} \quad \overline{\left\{ \begin{matrix} i \\ p-j \end{matrix} \right\}} = 0.$$

Indeed, the point is between the two cells, inside the zero hole; hence the relation is true.

(iii) $p - i \leq j < p$; hence $p - j \leq i$.

The point is in the free part of the right hand cell (on the negative side), its coordinates relative to the head being $(i, -(p - j))$.

Now by similarity and symmetry

$$\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}} = r \overline{\left\{ \begin{matrix} i \\ p-j \end{matrix} \right\}}$$

and, since $j \geq p - i > i$,

$$\overline{\left\{ \begin{matrix} i \\ j \end{matrix} \right\}} = 0.$$

The relation is verified again.

For $i > \frac{p-1}{2}$ we proceed by induction. We note first that the relation holds also for the axes of the two cells, namely, $j = 0$ and $j = p$, by reasoning as in cases (i) and (iii). For a point $P(n, k)$ in the lower half of the rectangle, we have the recursion formula (12). Since we may assume that $0 < j < p$, since the cases $j = 0$ and $j = p$ are settled, we assume that

(19) is valid for all three terms $\left\{ \begin{smallmatrix} n-1 \\ k-1 \end{smallmatrix} \right\}$, $\left\{ \begin{smallmatrix} n-1 \\ k \end{smallmatrix} \right\}$ and $\left\{ \begin{smallmatrix} n-1 \\ k+1 \end{smallmatrix} \right\}$. From (12) we have

$$\begin{aligned} \overline{\left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}} &= \left(\ell \overline{\left\{ \begin{smallmatrix} i-1 \\ j-1 \end{smallmatrix} \right\}} + r \overline{\left\{ \begin{smallmatrix} i-1 \\ p-(j-1) \end{smallmatrix} \right\}} \right) + \left(\ell \overline{\left\{ \begin{smallmatrix} i-1 \\ j \end{smallmatrix} \right\}} + r \overline{\left\{ \begin{smallmatrix} i-1 \\ p-j \end{smallmatrix} \right\}} \right) \\ &\quad + \left(\ell \overline{\left\{ \begin{smallmatrix} i-1 \\ j+1 \end{smallmatrix} \right\}} + r \overline{\left\{ \begin{smallmatrix} i-1 \\ p-(j+1) \end{smallmatrix} \right\}} \right). \end{aligned} \quad (20)$$

Applying the recursion formula for the principal cell, we have

$$\overline{\left\{ \begin{smallmatrix} i \\ j \end{smallmatrix} \right\}} = \overline{\left\{ \begin{smallmatrix} i-1 \\ j-1 \end{smallmatrix} \right\}} + \overline{\left\{ \begin{smallmatrix} i-1 \\ j \end{smallmatrix} \right\}} + \overline{\left\{ \begin{smallmatrix} i-1 \\ j+1 \end{smallmatrix} \right\}}.$$

The corresponding recursion holds for $\overline{\left\{ \begin{smallmatrix} i \\ p-j \end{smallmatrix} \right\}}$; hence by rearranging the terms in (20), we verify the relation (19). \square

Alternatively, (19) may be proved by the convolution identity (15):

$$\begin{aligned} \overline{\left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}} &= \sum_s \overline{\left\{ \begin{smallmatrix} ap \\ s \end{smallmatrix} \right\}} + \overline{\left\{ \begin{smallmatrix} i \\ k-s \end{smallmatrix} \right\}} \\ &= \overline{\left\{ \begin{smallmatrix} ap \\ bp \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} i \\ j \end{smallmatrix} \right\}} + \overline{\left\{ \begin{smallmatrix} ap \\ bp+p \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} i \\ j-p \end{smallmatrix} \right\}} \\ &= \ell \overline{\left\{ \begin{smallmatrix} i \\ j \end{smallmatrix} \right\}} + r \overline{\left\{ \begin{smallmatrix} i \\ j-p \end{smallmatrix} \right\}}. \end{aligned}$$

Note 1. For all i in the region

$$j < p - i \Rightarrow \overline{\left\{ \begin{smallmatrix} i \\ p-j \end{smallmatrix} \right\}} = 0$$

and

$$j > i \Rightarrow \overline{\left\{ \begin{smallmatrix} i \\ j \end{smallmatrix} \right\}} = 0;$$

hence the region of overlap is defined by $p - i \leq j \leq i$ (this implies that $i > \frac{p}{2}$).

Note 2. As remarked earlier, the principal cell may have zero entries and so some cells may have zero heads. Trivially, the theorem is still valid for these cases, including the case when the left cell is "extreme" and the right cell is out of the array (hence all its entries are zero).

As in binomial arrays, the last row of the principal cell has a simple structure. For $p = 2$ and $p = 3$, the situations are trivial, the $(p - 1)$ -st rows being 1 1 and 1 2 0 2 1, respectively.

Let $p > 3$. Now

$$\overline{\left\{ \begin{smallmatrix} p \\ -p \end{smallmatrix} \right\}} = \overline{\left\{ \begin{smallmatrix} p \\ 0 \end{smallmatrix} \right\}} = \overline{\left\{ \begin{smallmatrix} p \\ p \end{smallmatrix} \right\}} = 1 \quad \text{and} \quad \overline{\left\{ \begin{smallmatrix} p \\ k \end{smallmatrix} \right\}} = 0 \quad \text{for } p \neq k.$$

Thus, for $0 < i < p - 2$,

$$\overline{\left\{ \begin{matrix} p-1 \\ i-1 \end{matrix} \right\}} + \overline{\left\{ \begin{matrix} p-1 \\ i \end{matrix} \right\}} + \overline{\left\{ \begin{matrix} p-1 \\ i+1 \end{matrix} \right\}} = \overline{\left\{ \begin{matrix} p-1 \\ i \end{matrix} \right\}} + \overline{\left\{ \begin{matrix} p-1 \\ i+1 \end{matrix} \right\}} + \overline{\left\{ \begin{matrix} p-1 \\ i+2 \end{matrix} \right\}} = 0;$$

hence

$$\overline{\left\{ \begin{matrix} p-1 \\ i+2 \end{matrix} \right\}} = \overline{\left\{ \begin{matrix} p-1 \\ i-1 \end{matrix} \right\}}$$

in the above range. Thus $\overline{\left\{ \begin{matrix} p-1 \\ i \end{matrix} \right\}}$ is periodic with period = 3. Furthermore

$$\overline{\left\{ \begin{matrix} p \\ 0 \end{matrix} \right\}} = \overline{\left\{ \begin{matrix} p-1 \\ -1 \end{matrix} \right\}} + \overline{\left\{ \begin{matrix} p-1 \\ 0 \end{matrix} \right\}} + \overline{\left\{ \begin{matrix} p-1 \\ 1 \end{matrix} \right\}} = 1$$

or

$$\overline{\left\{ \begin{matrix} p-1 \\ 0 \end{matrix} \right\}} + 2\overline{\left\{ \begin{matrix} p-1 \\ 1 \end{matrix} \right\}} = 1. \quad (21)$$

If $p \equiv 1 \pmod{3}$, then by periodicity

$$\overline{\left\{ \begin{matrix} p-1 \\ 0 \end{matrix} \right\}} = \overline{\left\{ \begin{matrix} p-1 \\ p-1 \end{matrix} \right\}} = 1$$

and from (21)

$$\overline{\left\{ \begin{matrix} p-1 \\ 1 \end{matrix} \right\}} = 0;$$

so for $i \geq 0$ the entries form the sequence 1 0 -1 ... 1. If $p \equiv -1 \pmod{3}$ then

$$\overline{\left\{ \begin{matrix} p-1 \\ 1 \end{matrix} \right\}} = \overline{\left\{ \begin{matrix} p-1 \\ p-1 \end{matrix} \right\}} = 1$$

and from (21)

$$\overline{\left\{ \begin{matrix} p-1 \\ 0 \end{matrix} \right\}} = -1.$$

In this case the entries for $i \geq 0$ are -1 1 0 -1 1 0 ... 1.

It follows from this and (19) that the entries in row $p - 1$ of Figure 7 are

$$\ell, r, -(\ell + r), \ell, r, -(\ell + r), \dots, \ell, r, \quad \text{if } p \equiv 1 \pmod{3}$$

and

$$-\ell, \ell + r, -r, -\ell, \dots, \ell + r, \dots, \ell + r, -r, \quad \text{if } p \equiv -1 \pmod{3}.$$

The algorithm to find $\overline{\left\{ \begin{matrix} n \\ k \end{matrix} \right\}}$ in the general case can now be developed as follows.

Let

$$n = a_m p^m + a_{m-1} p^{m-1} + \dots + a_0$$

$$k = b_m p^m + b_{m-1} p^{m-1} + \dots + b_0$$

be the expansions to base p of n and k , where $0 \leq a_s, b_s < p$ for $0 \leq s \leq m$ and $a_m > 0$, $b_m \leq a_m$.

Let

$$A_1 = \sum_{t=1}^m a_t p^{t-1}, \quad B_1 = \sum_{t=1}^m b_t p^{t-1}.$$

Then

$$n = A_1 p + a_0, \quad k = B_1 p + b_0.$$

The heads of the cells related to the entry $\overline{\left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}}$ (referring to Figure 7) are

$$l = \overline{\left\{ \begin{smallmatrix} A_1 \\ B_1 \end{smallmatrix} \right\}} \quad \text{and} \quad r = \overline{\left\{ \begin{smallmatrix} A_1 \\ B_1 + 1 \end{smallmatrix} \right\}}$$

by (18). Hence, by (19),

$$\overline{\left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}} = \overline{\left\{ \begin{smallmatrix} A_1 \\ B_1 \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} a_0 \\ b_0 \end{smallmatrix} \right\}} + \overline{\left\{ \begin{smallmatrix} A_1 \\ B_1 + 1 \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} a_0 \\ p - b_0 \end{smallmatrix} \right\}}.$$

Thus, if $\overline{\left\{ \begin{smallmatrix} A_1 \\ B_1 \end{smallmatrix} \right\}}$ and $\overline{\left\{ \begin{smallmatrix} A_1 \\ B_1 + 1 \end{smallmatrix} \right\}}$ are known, then $\overline{\left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}}$ can be found.

More generally, let

$$A_s = \sum_{t=s}^m a_t p^{t-s} \quad \text{and} \quad B_s = \sum_{t=s}^m b_t p^{t-s}.$$

Then $\overline{\left\{ \begin{smallmatrix} A_{s+1} \\ B_{s+1} \end{smallmatrix} \right\}}$ and $\overline{\left\{ \begin{smallmatrix} A_{s+1} \\ B_{s+1} + 1 \end{smallmatrix} \right\}}$ determine $\overline{\left\{ \begin{smallmatrix} A_s \\ B_s \end{smallmatrix} \right\}}$ where $A_s = A_{s+1} p + a_s$ and $B_s = B_{s+1} p + b_s$. So

$$\overline{\left\{ \begin{smallmatrix} A_s \\ B_s \end{smallmatrix} \right\}} = \overline{\left\{ \begin{smallmatrix} A_{s+1} \\ B_{s+1} \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} a_s \\ b_s \end{smallmatrix} \right\}} + \overline{\left\{ \begin{smallmatrix} A_{s+1} \\ B_{s+1} + 1 \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} a_s \\ p - b_s \end{smallmatrix} \right\}} \quad (22)$$

and $\overline{\left\{ \begin{smallmatrix} A_s \\ B_s + 1 \end{smallmatrix} \right\}}$ is determined similarly. We note that

$$A_m = a_m, \quad B_m = b_m, \quad B_m + 1 = b_m + 1.$$

Hence, beginning with this, we have

$$\begin{aligned} \overline{\left\{ \begin{smallmatrix} A_{m-1} \\ B_{m-1} \end{smallmatrix} \right\}} &= \overline{\left\{ \begin{smallmatrix} a_m \\ b_m \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} a_{m-1} \\ b_{m-1} \end{smallmatrix} \right\}} + \overline{\left\{ \begin{smallmatrix} a_m \\ b_m + 1 \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} a_{m-1} \\ p - b_{m-1} \end{smallmatrix} \right\}} \\ \overline{\left\{ \begin{smallmatrix} A_{m-1} \\ B_{m-1} + 1 \end{smallmatrix} \right\}} &= \overline{\left\{ \begin{smallmatrix} a_m \\ b_m \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} a_{m-1} \\ b_{m-1} + 1 \end{smallmatrix} \right\}} + \overline{\left\{ \begin{smallmatrix} a_m \\ b_m + 1 \end{smallmatrix} \right\}} \overline{\left\{ \begin{smallmatrix} a_{m-1} \\ p - (b_{m-1} + 1) \end{smallmatrix} \right\}} \end{aligned}$$

Proceeding step by step, using (22), we evaluate $\overline{\left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}}$. The algorithm has been programmed, and gives fast results for large values of n and k . There is a closed formula consisting of sums of products (most of which are zero in practice), which we do not write down.

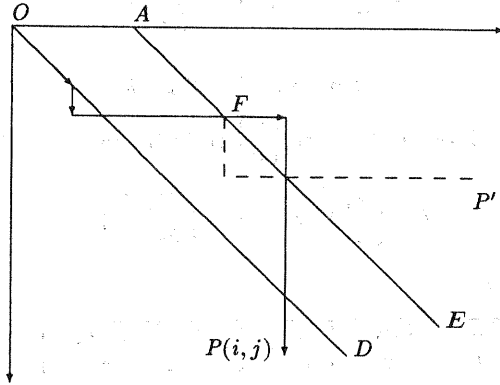


Figure 8

The algorithm simplifies to the familiar Lucas formula

$$\overline{\begin{Bmatrix} n \\ k \end{Bmatrix}} = \overline{\begin{Bmatrix} a_m \\ b_m \end{Bmatrix}} \overline{\begin{Bmatrix} a_{m-1} \\ b_{m-1} \end{Bmatrix}} \cdots \overline{\begin{Bmatrix} a_0 \\ b_0 \end{Bmatrix}}$$

if and only if the second term in (19) is 0; this means, geometrically, that the coordinates (a_s, b_s) determine a point in the free part of the cell, or in the zero-hole. This happens if and only if for all s in the expansions, $a_s + b_s < p$.

4. RESTRICTED PATHS

In both cases discussed in this paper, further restraints may be placed on the king's route. When progress is allowed only in the positive direction (Section 2) one possible restraint proposed in [1] is to cut the chessboard diagonally, that is, to allow only positions (i, j) where $j \leq i$. We denote by $u(i, j)$ the number of possible paths from $(0, 0)$ to (i, j) where each step, vertical, horizontal or diagonal is positively directed and $y \leq x$ for each point (x, y) along the path. Tables (5) and (6) are computer outputs showing arrays of $u(i, j) \bmod 3$ and $u(i, j) \bmod 7$ respectively.

Let $\bar{u}(i, j) = u(i, j) \pmod{p}$. It is easy to evaluate $\bar{u}(i, j)$ for all values of (i, j) , but the array has not the cell cluster structure in the sense that the array of the $\bar{u}(i, j)$ has. It has, however, some outstanding regularity features (inherited from the $\bar{u}(i, j)$ array).

A formula for the $u(i, j)$ function will be found first. The well known technique of *reflection* is used here.

On Figure 8, OD represents the line bisecting diagonally the chessboard, and the line AE is drawn parallel to it, originating at $A(0, 1)$.

The figure represents an "illegal" path, getting to (i, j) through some points in the "forbidden" half of the board. The point where the path first meets the line AE is F .

Leaving the section from 0 to F unchanged, reflect the section from F to P about AE . The reflection is shown by the dotted line ending at P' . The coordinates of P' are $(j-1, i+1)$. The reflection determines a one-to-one map of any illegitimate path to (i, j) to another path to $(j-1, i+1)$. Thus $u(i, j)$ is obtained by eliminating the illegal routes, and so

$$u(i, j) = w(i, j) - w(j-1, i+1). \quad (23)$$

Thus the problem of determining $\bar{u}(i, j)$ is reduced to finding $\bar{w}(i, j)$ and $\bar{w}(i+1, j-1)$.

Two noticeable features of the $\bar{u}(i, j)$ array are:

- (i) the apparent preservation of the zero-holes of the $w(i, j)$ array,
- (ii) the string of zeros running along diagonally at equal intervals.

Interpretation of (i) and (ii).

(i) From (23) and the symmetry of $\bar{w}(i, j)$ array it follows that

$$\bar{u}(i, j) = \bar{w}(i, j) - \bar{w}(i+1, j-1). \quad (24)$$

The zero-holes of the $\bar{w}(i, j)$ array represent square shaped cells or clusters containing $p \times p$, $p^2 \times p^2$, ... zero-entries. Exclude the first column and last row of the zero-hole. Then, for any entry (i, j) of the remaining array, $\bar{w}(i, j) = \bar{w}(i+1, j-1) = 0$; hence $\bar{u}(i, j) = 0$. Thus the $\bar{u}(i, j)$ array contains zero-holes beginning in rows divisible by p , p^2 , ...; these holes are generally square shaped and of size $(p-1) \times (p-1)$, $(p^2-1) \times (p^2-1)$, At the cut-off edge, the squares are cut to triangular shape.

(ii) Consider the point (d, d) on the diagonal of some cell in the $\bar{w}(i, j)$ array. By the symmetry of the cell, $\bar{w}(d, d+1) = \bar{w}(d+1, d)$, as long as both $(d, d+1)$ and $(d+1, d)$ belong to the same cell. (This is not the case when (d, d) is in the last row.)

Hence by (24), $\bar{u}(d, d+1) = \bar{w}(d, d+1) - \bar{w}(d+1, d) = 0$, provided that $p \nmid (d+1)$.

This accounts for the diagonal strings of zeros, cut at points where i is a multiple of p .

Next we deal with restrictions on the minimal route (Section 3). Outputs are shown on tables (7) and (8), of the positive half of the array mod 3 and mod 7. Here the restriction consists of cutting away the negative part of the chessboard, while still allowing diagonal steps in the negative direction, as long as the whole path is confined to the positive side.

Let AD , shown on Figure 9, be the line $y+1=0$, and $OFFP$ an illegal path, with F the first intersection of this path with AD , and the path FP' the reflection of the FP section in AD . The coordinates of P' are $(n, -(k+2))$.

Denote by $\left[\begin{smallmatrix} n \\ k \end{smallmatrix} \right]$ the number of restricted paths between 0 and P . All illegal paths ending at P are bijectively mapped to P' . Excluding these, we obtain

$$\left[\begin{smallmatrix} n \\ k \end{smallmatrix} \right] = \left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\} - \left\{ \begin{smallmatrix} n \\ -(k+2) \end{smallmatrix} \right\} = \left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\} - \left\{ \begin{smallmatrix} n \\ k+2 \end{smallmatrix} \right\} \quad (25)$$

Thus

$$\left[\begin{smallmatrix} n \\ k \end{smallmatrix} \right] = \left\{ \begin{smallmatrix} n-1 \\ k-1 \end{smallmatrix} \right\} + \left\{ \begin{smallmatrix} n-1 \\ k \end{smallmatrix} \right\} + \left\{ \begin{smallmatrix} n-1 \\ k+1 \end{smallmatrix} \right\} - \left(\left\{ \begin{smallmatrix} n-1 \\ k+1 \end{smallmatrix} \right\} + \left\{ \begin{smallmatrix} n-1 \\ k+2 \end{smallmatrix} \right\} + \left\{ \begin{smallmatrix} n-1 \\ k+3 \end{smallmatrix} \right\} \right)$$

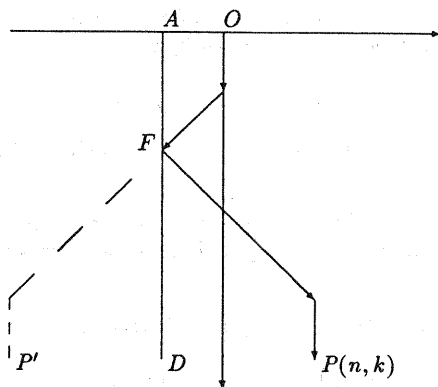


Figure 9

$$= \begin{bmatrix} n-1 \\ k-1 \end{bmatrix} + \begin{bmatrix} n-1 \\ k \end{bmatrix} + \begin{bmatrix} n-1 \\ k+1 \end{bmatrix} \quad (k \geq 0)$$

with

$$\begin{bmatrix} 0 \\ 0 \end{bmatrix} = 1.$$

Let

$$\overline{\begin{bmatrix} n \\ k \end{bmatrix}} = \begin{bmatrix} n \\ k \end{bmatrix} \pmod{p}.$$

By (25),

$$\overline{\begin{bmatrix} n \\ k \end{bmatrix}} = \overline{\begin{bmatrix} n \\ k \end{bmatrix}} - \overline{\begin{bmatrix} n \\ k+2 \end{bmatrix}} \quad (26)$$

From (26) it follows that

$$(i) \quad \overline{\begin{bmatrix} np \\ k \end{bmatrix}} = 0 \text{ unless } p|k \text{ or } p|(k+2)$$

$$(ii) \quad \overline{\begin{bmatrix} np \\ kp \end{bmatrix}} = \overline{\begin{bmatrix} np \\ kp \end{bmatrix}} - \overline{\begin{bmatrix} np \\ kp+2 \end{bmatrix}} = \overline{\begin{bmatrix} np \\ kp \end{bmatrix}} = \overline{\begin{bmatrix} n \\ k \end{bmatrix}} \quad (k \geq 0).$$

Note that, in general, $\overline{\begin{bmatrix} np \\ kp \end{bmatrix}} \neq \overline{\begin{bmatrix} n \\ k \end{bmatrix}}$, and this means that the cell cluster structure of the $\begin{bmatrix} n \\ k \end{bmatrix}$ array is not inherited.

However, an interesting feature of the array is the regularly spaced vertical columns of zeros, interrupted in the rows just above the rows where $n \equiv 0 \pmod{p}$. Consider

$$\overline{\begin{bmatrix} A_1p + a_0 \\ B_1p + b_0 \end{bmatrix}}, \text{ where } b_0 = p-1 \text{ and } 0 \leq a_0 < p-1.$$

Then, by (26),

$$\overline{\begin{bmatrix} A_1p + a_0 \\ B_1p + p - 1 \end{bmatrix}} = \overline{\left\{ \begin{array}{c} A_1p + a_0 \\ (B_1 + 1)p - 1 \end{array} \right\}} - \overline{\left\{ \begin{array}{c} A_1p + a_0 \\ (B_1 + 1)p + 1 \end{array} \right\}}$$

The two terms on the right hand side represent entries in a cell where the head is

$$\left\{ \begin{array}{c} A_1p \\ (B_1 + 1)p \end{array} \right\}.$$

If $a_0 < p - 1$, the entries are in the *free* part of the cell which is symmetrical; hence the terms are equal. Thus under the above conditions

$$\overline{\begin{bmatrix} A, p + a_0 \\ B, p + b_0 \end{bmatrix}} = 0.$$

REFERENCES

1. M. Razpet, "Divisibility properties of some number arrays", to appear.
2. M. Sved, "Geometry of Combinatorial Arithmetic", *Ars Combinatoria* 21-A (1986), 271-298.
3. M. Sved, "Divisibility—With Visibility", *The Mathematical Intelligencer*, 10(2) (1988), 56-64.

KING-WALK MODULO 3
 A, B, C, D ARE 1 1 0

0	
1	1
2	111
3	12.21
4	1.1.1.1
5	1111111111
6	12.....21
7	1..2.....2..1
8	111222...222111
9	12.12.12.21.21.21.21
10	1.....1.....1
11	111.....111.....111
12	12.21...12.21...12.21
13	1..1..1..1..1..1..1..1..1..1
14	11111111111111111111111111111111
15	12.....21
16	1..2.....2..1
17	111222...222111
18	12.12.12.....21.21.21
19	1.....2.....1
20	111...222.....222.....111
21	12.21...21.12.....21.12...12.21
22	1..1..1..2..2..2.....2..2..2..1..1..1
23	111111111222222222.....222222221111111111
24	12.....12.....12.....21.....21.....21
25	1..2.....1..2.....1..2.....2..1.....2..1.....2..1
26	111222...111222...111222...222111...222111...222111
27	12.12.12.12.12.12.12.12.12.21.21.21.21.21.21.21.21
28	1.....1.....1
29	111.....111.....111
30	12.21.....12.21.....12.21
31	1..1..1.....1..1..1.....1..1..1
32	111111111.....111111111.....111111111
33	12.....21.....12.....21.....12.....21
34	1..2.....2..1.....1..2.....2..1.....1..2.....2..1
35	111222...222111.....111222...222111.....111222...222111
36	12.12.12.21.21.21.....12.12.12.21.21.21.....12.12.12.21.21.21
37	1.....1.....1.....1.....1.....1.....1.....1
38	111.....111.....111.....111.....111.....111.....111.....111
39	12.21...12.21...12.21...12.21...12.21...12.21...12.21...12.21
40	1..1
41	11
42	12.....21
43	1..2.....2..1
44	111222...222111.....2..1
	12.12.12.....222111
21.21.21

TABLE 3

RESTRICTED WALK MODULO 3
 A, B, C, D ARE 1 1 1 0

0 1
 1 12
 2 11.
 3 1.11
 4 12.2.
 5 11.211
 6 1.1..2.
 7 12.11.22
 8 11.1.1.1.
 9 1.12.12.11
 10 12..2..2.2..
 11 11..211.21..
 12 1.11..2.22...
 13 12.2..21.1....
 14 11.211..12.....
 15 1.1..2..11.....
 16 12.11.22.2.....
 17 11.1.1.1.211111111
 18 1.12.12.1..21.21.2.
 19 12..2..2.11..1..1.22
 20 11..211.2.2..122.1.1.
 21 1.11..2.21.22..1.12.11
 22 12.2..21..1.1..12..2.2.
 23 11.211..112.122..221.211
 24 1.1..2..1..2..1..2..1..2.
 25 12.11.22.11.22.11.22.11.22
 26 11.1.1.1.1.1.1.1.1.1.1.1.
 27 1.12.12.12.12.12.12.12.11
 28 12..2..2..2..2..2..2..2..2.
 29 11..211.221..211.221..211.21..
 30 1.11..2.2..11..2.2..11..2.22...
 31 12.2..21.22.2..21.22.2..21.1....
 32 11.211..1.1.211..1.1.211..12.....
 33 1.1..2..12.1..2..12.1..2..11.....
 34 12.11.22..2.11.22..2.11.22.2.....
 35 11.1.1.1..212121211.2.2.2.21.....
 36 1.12.12.11.....2.21.21.22.....
 37 12..2..2.2.....21..1..1.1.....
 38 11..211.21.....22..122.12.....
 39 1.11..2.22.....2.22..1.11.....
 40 12.2..21.1.....21.1..12.2.....
 41 11.211..12.....22.122..21.....
 42 1.1..2..11.....2.2..1..22.....
 43 12.11.22.2.....21.22.11.1.....
 44 11.1.1.1.211111111..1.1.1.12.....

TABLE 5

