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**Abstract.** Let  $c(T_n)$  denote the number of 3-cycles in the tournament  $T_n$  and let  $u(T_n)$  denote the number of nodes  $i$  in  $T_n$  such that each arc oriented towards  $i$  belongs to at least one 3-cycle. We determine the minimum value of  $c(T_n)$  when  $u(T_n) = n$  and the maximum value of  $c(T_n)$  when  $u(T_n) = 3$ .

**1. Introduction.** A tournament  $T_n$  consists of a set of  $n$  nodes  $1, 2, \dots, n$  such that each pair of distinct nodes  $i$  and  $j$  is joined by exactly one of the arcs  $\overrightarrow{ij}$  or  $\overrightarrow{ji}$ . If the arc  $\overrightarrow{ij}$  is in  $T_n$  we say that  $i$  beats  $j$  or that  $j$  loses to  $i$  and write  $i \rightarrow j$ . If each node of a subtournament  $A$  beats each node of a subtournament  $B$  we write  $A \rightarrow B$ . For definitions not given here or for additional material on tournaments, see [12] or [15].

Node  $i$  is said to cover node  $j$  if node  $\overrightarrow{i}$  beats every node that node  $j$  beats or, equivalently, if  $i \rightarrow j$  and the arc  $\overrightarrow{ij}$  belongs to no 3-cycle. It is not difficult to see that the covering relation thus defined is transitive [10; p. 72]. So, as pointed out in [10], every finite tournament has at least one uncovered node (a result originally proved by another argument in [8; p. 148]). In fact, every strong tournament  $T_n$  with  $n \geq 3$  nodes has at least three uncovered nodes (cf. [17], [11], [15; p. 178], [9] or [10]). We observe that node  $i$  is an uncovered node if and only if every arc oriented towards  $i$  belongs to at least one 3-cycle or, equivalently, if for any other node  $j$  there exists a path from  $i$  to  $j$  of length at most two; nodes with this property have been called *kings* in several recent papers (cf. [9], [2], [18], [7] or [4]).

Miller [10] has shown that if the tournament  $T_n$  represents majority preferences between a set of  $n$  proposals, then various voting procedures will always select a proposal from the set of uncovered nodes of  $T_n$  (see also [16], [1], [14], or [5] for additional material involving uncovered nodes in this context). Let  $u(T_n)$  denote the number of uncovered nodes in the tournament  $T_n$ . Miller [10; p. 78] remarked that the size of  $u(T_n)$ , in a strong tournament  $T_n$ , "depends largely on the degree of intransitivity within  $[T_n]$ , which in turn may be measured by the proportion of all triples of [nodes in  $T_n$ ] that are cyclic  $\dots$ . With some

complications, the pattern is this: as intransitivity  $\dots$  declines,  $[u(T_n)]$  approaches or equals 3; as it increases,  $[u(T_n)]$  approaches or equals  $[n]$ ." Our main object here is to determine the minimum number of 3-cycles possible in a tournament  $T_n$  when  $u(T_n) = n$  and the maximum number of 3-cycles possible when  $u(T_n) = 3$ .

**2. Preliminary Remarks.** The *score*  $s_i$  of a node  $i$  in a tournament  $T_n$  is the number of nodes beaten by  $i$  and the *score sequence* of  $T_n$  is the sequence  $\bar{s} = (s_1, s_2, \dots, s_n)$ ; the sum of the scores is clearly  $n(n-1)/2$ . We let  $R_n$  denote any *regular* tournament with  $n$  nodes, that is, a tournament in which  $s_i = (n-1)/2$  for all nodes if  $n$  is odd, or  $s_i = n/2 - 1$  for half the nodes and  $n/2$  for the other half if  $n$  is even.

It is well-known and easy to see that the number  $c(T_n)$  of 3-cycles in a tournament  $T_n$  with score sequence  $\bar{s}$  is given by the formula

$$(2.1) \quad c(T_n) = \binom{n}{3} - \sum_1^n \binom{s_i}{2};$$

from this it follows readily that

$$(2.2) \quad c(T_n) \leq \gamma(n) := \begin{cases} (n^3 - n)/24 & \text{if } n \text{ is odd} \\ (n^3 - 4n)/24 & \text{if } n \text{ is even,} \end{cases}$$

with equality holding if and only if  $T_n$  is a regular tournament  $R_n$  (see, e.g., [12; p. 9] or [15; p. 186]). Landau [8] showed that every node of maximum score in a tournament  $T_n$  is an uncovered node. Miller [10; p. 80] pointed out that this and result (2.2) imply that if  $c(T_n) = \gamma(n)$  then  $u(T_n) = n$  if  $n$  is odd, and  $n/2 \leq u(T_n) \leq n$  if  $n$  is even. When  $n = 2m$  the upper bound here can be realized for all  $m \geq 3$  and the lower bound can be realized if and only if  $m$  is odd; if  $n = 4m$  then tournaments  $T_n$  such that  $c(T_n) = \gamma(n)$  and  $u(T_n) = n/2 + 1$  exist for all  $m \geq 1$ . In Section 4 we show that if  $u(T_n) = n$  then the minimum possible value of  $c(T_n)$  is a quadratic in  $n$ , with leading term  $n^2/4$ , and we determine the minimal tournaments.

In the other direction, if  $T_n$  is a strong tournament with  $n \geq 3$  nodes, then  $c(T_n) \geq n-2$  [6; p. 306]. Burzio and Demaria [3] (see also [13]) have recently characterized the tournaments for which equality holds here and, as it turns out, all these tournaments have exactly three uncovered nodes. In Section 5 we show that if  $u(T_n) = 3$ , then the maximum possible value of  $c(T_n)$  is a cubic in  $n$ , whose first two terms are  $(n^3 - n^2)/24$ , and we characterize the maximal tournaments.

**3. An Inequality.** Let  $\bar{s} = (s_1, s_2, \dots, s_n)$  be the score sequence of a tournament  $T_n$  such that  $u(T_n) = n$ . We may suppose the nodes are labelled so that

$$(3.1) \quad s_1 \leq s_2 \leq \dots \leq s_n.$$

Let  $S_j$  and  $e_j$  be such that  $S_j = s_1 + \dots + s_j = j(j-1)/2 + e_j$  for  $1 \leq j \leq n$ ; then  $e_j \geq 0$  since there are  $j(j-1)/2$  arcs joining any  $j$  nodes. Suppose that  $j + e_j \leq n - 1$  and that  $e_j \leq j - 1$  for some  $j$ . Then there would be a node  $p \in \{j+1, \dots, n\}$  that beats all nodes in  $\{1, \dots, j\}$  and a node  $q \in \{1, \dots, j\}$  that loses to all nodes in  $\{j+1, \dots, n\}$ ; but then node  $p$  would cover node  $q$ , contrary to the assumption that  $u(T_n) = n$ . Consequently, if  $u(T_n) = n$ , then

$$(3.2) \quad S_j \geq \binom{j}{2} + \min\{j, n-j\} = \binom{j+1}{2} + \min\{0, n-2j\}$$

for  $1 \leq j \leq n-1$ , and

$$(3.3) \quad S_n = \binom{n}{2}.$$

We now derive an inequality involving such sequences  $\bar{s}$  that we shall use in the next section to determine the minimum value of  $c(T_n)$  if  $u(T_n) = n$ .

LEMMA 1. Let  $\bar{s} = (s_1, \dots, s_n)$  denote a sequence of  $n \geq 5$  integers satisfying conditions (3.1)-(3.3). Let  $\bar{x} = (1, 2, \dots, m-1, m, m, m, m+1, \dots, n-2)$  if  $n = 2m+1 \geq 5$  and let  $\bar{y} = (1, 2, \dots, m-1, m, m, m, m, m+2, \dots, n-2)$  and  $\bar{z} = (1, 2, \dots, m-3, m-1, m-1, m-1, m-1, m, \dots, n-2)$  if  $n = 2m \geq 8$ ; finally, let  $\bar{y} = (1, 2, 3, 3, 3, 3)$  and  $\bar{z} = (2, 2, 2, 2, 3, 4)$  if  $n = 6$ . Then

$$(3.4) \quad f(\bar{s}) := \sum_1^n \binom{s_i}{2} \leq \binom{n-1}{3} + \begin{cases} (n-1)(n-3)/4 & \text{if } n = 2m+1 \\ n(n-4)/4 & \text{if } n = 2m, \end{cases}$$

with equality holding if and only if  $\bar{s} = \bar{x}$  when  $n$  is odd or  $\bar{s} = \bar{y}$  or  $\bar{z}$  when  $n$  is even.

PROOF: It is easy to verify that equality holds in (3.4) when  $\bar{s} = \bar{x}$ ,  $\bar{y}$ , or  $\bar{z}$ . Let us assume that  $\bar{s}$  is a maximal sequence, i.e., a sequence for which  $f(\bar{s})$  assumes its maximum value over the set of all sequences satisfying (3.1)-(3.3). Suppose there

exists a least integer  $k$  such that strict inequality holds in (3.2) when  $j = k$ . We assume, initially, that  $1 \leq k \leq m$  so that

$$(3.5) \quad S_k = \binom{k+1}{2} + \alpha$$

where  $\alpha \geq 1$ . Then  $s_1 = 1 + \alpha$  if  $k = 1$ ; and if  $k \geq 2$  then it follows from the definition of  $k$  and the relation  $S_j = S_{j-1} + s_j$  that  $s_j = j$  for  $1 \leq j \leq k-1$ , and that

$$(3.6) \quad s_k = k + \alpha > s_{k-1}.$$

Let  $h$  denote the largest integer such that  $s_{k+1} = \dots = s_{k+h}$ ; then

$$(3.7) \quad s_{k+h} < s_{k+h+1}$$

if  $k+h < n$ . We now show that - apart from one exceptional case -

$$(3.8) \quad S_{k+u} > \binom{k+u+1}{2} + \min \{0, n - 2k - 2u\}$$

for  $1 \leq u \leq h-1$ , assuming that  $h \geq 2$ .

We observe that if  $s_{k+1} = \sigma$ , then

$$S_{k+u} = S_k + u\sigma = \binom{k+1}{2} + \alpha + u\sigma,$$

so (3.8) holds if and only if

$$(3.9) \quad u\{\sigma - k - (u+1)/2\} + \alpha + \max \{0, 2k + 2u - n\} > 0.$$

It follows from (3.2), (3.5), and the definition of  $h$  that

$$\begin{aligned} w\sigma = S_{k+w} - S_k &\geq \binom{k+w+1}{2} + \min \{0, n - 2k - 2w\} - \binom{k+1}{2} - \alpha \\ &= w\{k + (w+1)/2\} - \alpha - \max \{0, 2k + 2w - n\} \end{aligned}$$

for  $1 \leq w \leq h$ . So, in particular,

$$(3.10) \quad \sigma \geq k + (h+1)/2 - \alpha/h - h^{-1} \cdot \max \{0, 2k + 2h - n\}.$$

Moreover, if  $m + 1 \leq k + h$  and  $v := m + 1 - k$ , then

$$(3.11) \quad \sigma \geq k + (v + 1)/2 - v^{-1} \cdot (2m + 2 - n + \alpha).$$

We now apply these estimates in (3.9), considering three cases separately.

**Case 1:**  $k + u < k + h \leq m$ . Let  $L$  denote the left hand side of inequality (3.9). In this case it follows from (3.10) that

$$L \geq u\{(h - u)/2 - \alpha/h\} + \alpha \geq u/2 + \alpha/h > 0,$$

as required.

**Case 2:**  $k + u \leq m < k + v = m + 1 \leq k + h$ . Notice that  $v \geq u + 1 \geq 2$  here. In this case it follows from (3.11) that

$$L \geq u\{(v - u)/2 - v^{-1} \cdot (2m + 2 - n + \alpha)\} + \alpha \geq u/2 + \alpha/v - u(2m + 2 - n)/v.$$

If  $n = 2m + 1$ , then

$$L \geq u/2 + \alpha/v - u/v \geq \alpha/v > 0,$$

as required. If  $n = 2m$  and  $v \geq 4$ , which is certainly the case if  $u \geq 3$ , then

$$L \geq u/2 + \alpha/v - 2u/v \geq \alpha/v > 0.$$

Moreover, it follows from the inequality  $\sigma = s_{k+1} \geq s_k = k + \alpha$  that

$$L \geq u\{\alpha - (u + 1)/2\} + \alpha = (u + 1)(\alpha - u/2),$$

so  $L > 0$  if  $u = 1$  or  $u = 2$  and  $\alpha \geq 2$ . Thus we find that  $L > 0$  here *except* when  $n = 2m$ ,  $\alpha = 1$ ,  $u = 2$ ,  $v = 3$ , and  $s_{k+1} = k + 1$ ; in this exceptional case  $k = m + 1 - v = m - 2$  and  $(s_1, \dots, s_{m+1}) = (1, 2, \dots, m - 3, m - 1, m - 1, m - 1, m - 1)$ .

**Case 3:**  $m + 1 \leq k + u < k + h$ . In this case it follows from (3.10) that

$$\begin{aligned} L &\geq u\{(h - u)/2 - (\alpha + 2k - n)/h\} + (\alpha + 2k - n) \\ &= h^{-1}(h - u)\{\alpha + 2u + 2k - n + u(h - 4)/2\} \\ &\geq h^{-1}(h - u)\{\alpha + 1 + u(h - 4)/2\}, \end{aligned}$$

so  $L$  is certainly positive if  $h \geq 4$ ; and if  $2 \leq h \leq 3$ , then

$$\begin{aligned} L &\geq h^{-1}(h-u)\{\alpha+1+(h-1)(h-4)/2\} \\ &= h^{-1}(h-u)\{\alpha+(h-2)(h-3)/2\} \geq \alpha/h > 0, \end{aligned}$$

as required. Thus it follows that inequalities (3.9) and (3.8) hold apart from the one exceptional case.

Suppose we are not in this exceptional case and let  $\bar{r} = (r_1, \dots, r_n)$  denote the integer sequence in which  $r_i = s_i$  except that  $r_k = s_k - 1$  and  $r_{k+h} = s_{k+h} + 1$ . Then  $\bar{r}$  clearly satisfies condition (3.3); and  $\bar{r}$  also satisfies conditions (3.1) and (3.2) in view of inequalities (3.6) and (3.7) and relations (3.5) and (3.8), respectively. But

$$\begin{aligned} (3.12) \quad f(\bar{r}) - f(\bar{s}) &= \binom{s_{k+h}+1}{2} - \binom{s_{k+h}}{2} + \binom{s_k-1}{2} - \binom{s_k}{2} \\ &= s_{k+h} - s_k + 1 \geq 1, \end{aligned}$$

contrary to the assumption that  $\bar{s}$  is a maximal sequence.

It follows, therefore, that if  $\bar{s}$  is a maximal sequence then either strict equality holds in (3.2) for  $1 \leq j \leq m$  and

$$(3.13) \quad (s_1, \dots, s_m) = (1, 2, \dots, m),$$

or  $\bar{s}$  involves the exceptional case encountered earlier and  $n = 2m$  and

$$(3.14) \quad (s_1, \dots, s_{m+1}) = (1, 2, \dots, m-3, m-1, m-1, m-1, m-1).$$

(This last sequence is to be interpreted as  $(2, 2, 2, 2)$  if  $m = 3$ .) And, similarly, it follows by duality that if  $\bar{s}$  is a maximal sequence, then either  $(s_{n-m}, \dots, s_n) = (n-m-1, n-m, \dots, n-2)$  — that is,

$$(3.15) \quad (s_{m+2}, \dots, s_n) = (m, m+1, \dots, n-2)$$

if  $n = 2m + 1$  or

$$(3.16) \quad (s_{m+1}, \dots, s_n) = (m-1, m, \dots, n-2)$$

if  $n = 2m$  — or  $n = 2m$  and

$$(3.17) \quad (s_{m-1}, \dots, s_n) = (m, m, m, m, m+2, \dots, n-2).$$

If  $n = 2m + 1$  and we combine (3.13) and (3.15), we find that  $s_{m+1}$  must equal  $m$ , in view of (3.1), so  $\bar{s} = \bar{x}$ . If  $n = 2m$  and we combine the only compatible alternatives, namely, (3.13) and (3.17) or (3.14) and (3.16), we find that  $\bar{s} = \bar{y}$  or  $\bar{z}$ . This suffices to complete the proof of the lemma.

**4. Minimal Tournaments  $T_n$  with  $u(T_n) = n$ .** If  $u(T_6) = 6$  for a tournament  $T_6$ , then  $T_6$  must have score sequence  $(2, 2, 2, 3, 3, 3)$ ; this follows from an argument that will be given later. There are five non-isomorphic tournaments with this score sequence (cf. [12; p. 95]) and of these only the following three have the property that  $u(T_6) = 6$ : (i) the tournament  $T_6$  consisting of two disjoint 3-cycles  $(A, B, C)$  and  $(c, b, a)$  such that  $(A, B, C) \rightarrow (c, b, a)$  except that  $a \rightarrow A$ ,  $b \rightarrow B$ , and  $c \rightarrow C$ ; (ii) the tournament  $T_6$  with nodes  $1, 2, \dots, 6$  in which  $j \rightarrow i$  if  $j > i$  except that  $1 \rightarrow 5$ ,  $1 \rightarrow 6$ ,  $2 \rightarrow 4$ , and  $4 \rightarrow 6$ ; and (iii) the dual of the tournament described in (ii). Let  $M_1$  denote the trivial tournament with just one node and let  $M_6$  denote any one of the three tournaments just described. More generally, if  $n = 3$  or  $5$  or  $n \geq 7$  let  $M_n$  denote any tournament obtained from any tournament  $M_{n-2}$  by adjoining two nodes  $p$  and  $q$  such that  $p \rightarrow q$ ,  $q \rightarrow M_{n-2}$ , and  $M_{n-2} \rightarrow p$ . It is not difficult to verify that  $u(M_n) = n$  for any such tournament  $M_n$ . We now show that among all tournaments  $T_n$  such that  $u(T_n) = n$  these are the minimal tournaments, that is, the tournaments with the minimum number of 3-cycles.

**THEOREM 1.** Let  $T_n$  be a tournament with  $n \neq 2, 4$  nodes such that  $u(T_n) = n$ . Then

$$c(T_n) \geq \begin{cases} (n-1)^2/4 & \text{if } n \text{ is odd} \\ (n^2 - 2n + 8)/4 & \text{if } n \text{ is even,} \end{cases}$$

with equality holding if and only if  $T_n$  is one of the tournaments  $M_n$ .

**PROOF:** We may suppose that  $n \geq 5$  since the result certainly holds when  $n = 1$  or  $3$ . And, as we saw earlier, we may suppose the score sequence  $\bar{s}$  of  $T_n$  satisfies conditions (3.1)-(3.3) so that, in particular,  $s_1 \geq 1$  and  $s_n \leq n - 2$ .

We consider first the case when  $2 \leq s_1 \leq \dots \leq s_n \leq n - 3$ . If  $n = 6$  then  $\bar{s} = (2, 2, 2, 3, 3, 3)$  so  $c(T_6) = 8$ , by (2.1), and  $T_6$  is one of the tournaments  $M_6$ , in view of the earlier observation. So we may now suppose that  $n = 5$  or  $n \geq 7$ . If  $n$  is odd then  $\bar{s}$  is certainly not the sequence  $\bar{x}$  described in

Lemma 1; then in this case it follows from (2.1) and Lemma 2 that

$$\begin{aligned} c(T_n) &= \binom{n}{3} - f(\bar{s}) \\ &> \binom{n}{3} - f(\bar{x}) \\ &= \binom{n}{3} - \binom{n-1}{3} - (n-1)(n-3)/4 = (n-1)^2/4, \end{aligned}$$

as required.

If  $n \geq 8$  is even then  $\bar{s}$  is neither of the sequences  $\bar{y}$  or  $\bar{z}$  described in Lemma 1, so  $f(\bar{s}) < f(\bar{y}) = f(\bar{z})$ . Now  $\bar{s}$  can be transformed into one of the sequences  $\bar{y}$  or  $\bar{z}$  by a series of exchanges each of which involves replacing two elements  $s_k$  and  $s_{k+h}$  by  $s_k - 1$  and  $s_{k+h} + 1$ , respectively, where  $k$  and  $h$  are as defined in the proof of Lemma 1. Each such exchange increases the value of the sum  $f(\bar{s})$  by  $s_{k+h} - s_k + 1 \geq 1$ . Thus it follows from (3.12) that  $f(\bar{s}) + 1 \leq f(\bar{y}) = f(\bar{z})$  with equality holding only if  $\bar{s}$  can be transformed into  $\bar{y}$  or  $\bar{z}$  by a *single* exchange that involves replacing two *equal* elements  $s_k$  and  $s_{k+h}$  by  $s_k - 1$  and  $s_{k+h} + 1$ . Now  $y_1 = z_1 = 1$  and  $y_n = z_n = n - 2$ ; hence, if  $\bar{s}$  can be so transformed, it must be that  $k = 1$ ,  $k + h = n$ ,  $s_1 = 2$ , and  $s_n = n - 3$ . But we are assuming that  $n \geq 8$  here, so  $s_1$  cannot equal  $s_n$  and, consequently,  $\bar{s}$  cannot be so transformed into  $\bar{y}$  or  $\bar{z}$  by a single exchange.

We conclude, therefore, that if  $n$  is even,  $n \geq 8$ , and  $2 \leq s_1 \leq s_n \leq n - 3$ , then  $f(\bar{s}) + 1 < f(\bar{y}) = f(\bar{z})$ . So in this case it follows from (2.1) and Lemma 1 that

$$\begin{aligned} c(T_n) &= \binom{n}{3} + f(\bar{s}) \\ &> 1 + \binom{n}{3} - f(\bar{y}) \\ &= 1 + \binom{n}{3} - \binom{n-1}{3} - n(n-4)/4 \\ &= (n^2 - 2n + 8)/4, \end{aligned}$$

as required.

It remains to consider the case when  $s_1 = 1$  or  $s_n = n - 2$ . If there is a node  $p$  of score 1 let  $q$  denote the node that loses to  $p$  and let  $T_{n-2}$  denote the subtournament determined by the remaining nodes, so that  $p \rightarrow q$



and  $T_{n-2} \rightarrow p$ . If node  $q$  lost to some node  $w$  of  $T_{n-2}$  then  $w$  would cover node  $p$ , contrary to our hypothesis; it follows, therefore, that  $q \rightarrow T_{n-2}$  so  $q$  has score  $n-2$ . Similarly, if we initially assume there is a node of score  $n-2$  we find that the node that beats this node has score 1. Thus if  $s_1 = 1$  or  $s_n = n-2$  there exist nodes  $p$  and  $q$  of score 1 and  $n-2$ , respectively, such that  $T_n$  has the structure described above; and, in this case, it follows readily that

$$(4.1) \quad c(T_n) = n - 2 + c(T_{n-2}).$$

The subtournament  $T_{n-2}$  is such that  $u(T_{n-2}) = n-2$ ; for if in  $T_{n-2}$  some node  $v$  covers some node  $w$ , then  $v$  clearly covers  $w$  in  $T_n$  as well, contrary to our hypothesis. We now observe that there is no tournament  $T_4$  such that  $u(T_4) = 4$  (since there is no tournament  $T_2$  such that  $u(T_2) = 2$ ). Consequently, the only tournaments  $T_6$  such that  $u(T_6) = 6$  are those with score sequence  $(2, 2, 2, 3, 3, 3)$  that were discussed earlier. So, in completing the argument for the case when  $s_1 = 1$  and  $s_n = n-2$ , we may assume that  $n = 5$  or  $n \geq 7$  and that the required result has already been proved when  $n$  is replaced by  $n-2$ . Thus it follows from (4.1), the fact that  $u(T_{n-2}) = n-2$ , and the induction hypothesis, that

$$c(T_n) \geq n - 2 + \begin{cases} (n-3)^2/4 = (n-1)^2/4 & \text{if } n \text{ is odd} \\ (n-2)(n-4)/4 + 2 = (n^2 - 2n + 8)/4 & \text{if } n \text{ is even,} \end{cases}$$

with equality holding if and only if  $T_{n-2}$  is one of the tournaments  $M_{n-2}$ ; that is, if and only if  $T_n$  is one of the tournaments  $M_n$ . This suffices to complete the proof of the theorem.

**5. Minimizing Certain Sums.** In the next section we shall make use of the following slight extension of a familiar result on the minimum value of a sum  $\sum_1^n \binom{w_i}{2}$ , subject to the condition that the  $w_i$ 's are non-negative integers having a fixed sum.

LEMMA 2. Let  $J, K, j$ , and  $k$  be given positive integers such that

$$(5.1) \quad \lceil J/j \rceil \leq \lfloor K/k \rfloor.$$

For any integer  $D$  such that  $0 \leq D \leq J$ , let  $h(D)$  denote the minimum value of the sum

$$f(\bar{w}) = \sum_1^{j+k} \binom{w_i}{2}$$

over all sequences  $\bar{w} = (w_1, \dots, w_{j+k})$  of  $j+k$  non-negative integers such that

$$(5.2) \quad w_1 + \dots + w_j = J - D$$

and

$$(5.3) \quad w_{j+1} + \dots + w_{j+k} = K + D.$$

Then  $h(D)$  is a strictly increasing function of  $D$ ; furthermore, if  $\bar{w}$  satisfies conditions (5.2) and (5.3), then  $f(\bar{w}) = h(D)$  if and only if the integers  $w_1, \dots, w_j$  are as nearly equal as possible and the integers  $w_{j+1}, \dots, w_{j+k}$  are as nearly equal as possible.

PROOF: If  $1 \leq D \leq J$  let  $\bar{w}$  be a sequence that satisfies (5.2) and (5.3) and is such that  $f(\bar{w}) = h(D)$ . We may suppose that

$$w_1 \leq [(J - D)/j] \leq [J/j] - 1 \quad \text{and} \quad w_{j+k} \geq [(K + D)/k] \geq [K/k] + 1$$

so that  $w_{j+k} \geq w_1 + 2$ , by (5.1). Let  $\bar{w}'$  denote the sequence that differs from  $\bar{w}$  only in that  $w'_1 = w_1 + 1$  and  $w'_{j+k} = w_{j+k} - 1$ . Then  $\bar{w}'$  satisfies (5.2) and (5.3) with  $D$  replaced by  $D - 1$ . Moreover,

$$h(D) = f(\bar{w}) = f(\bar{w}') + w_{j+k} - w_1 - 1 \geq f(\bar{w}') + 1 \geq h(D - 1) + 1.$$

This proves the first part of the required conclusion; and the last part follows readily upon considering the subsequences  $(w_1, \dots, w_j)$  and  $(w_{j+1}, \dots, w_{j+k})$  separately.

**6. Maximal Tournaments  $T_n$  with  $u(T_n) = 3$ .** If  $n \geq 3$  let  $T_X, T_Y$ , and  $T_Z$  denote (possibly empty) tournaments with  $X, Y$  and  $Z$  nodes such that  $X + Y + Z = n - 3$ . Let  $Q = Q(T_X, T_Y, T_Z)$  denote the tournament consisting of disjoint copies of  $T_X, T_Y$ , and  $T_Z$  plus three additional nodes

$x, y,$  and  $z$  such that  $x \rightarrow T_X, y \rightarrow T_Y, z \rightarrow T_Z, \{x\} \cup T_X \rightarrow \{y\} \cup T_Y,$   
 $\{y\} \cup T_Y \rightarrow \{z\} \cup T_Z,$  and  $\{z\} \cup T_Z \rightarrow \{x\} \cup T_X.$  The only uncovered nodes in  
any such tournament  $Q$  are the nodes  $x, y,$  and  $z.$  We now show that among  
the class of tournaments  $T_n$  such that  $u(T_n) = 3,$  the maximal tournaments  
– that is, the tournaments with the maximum number of 3-cycles – are a certain  
subset of these tournaments  $Q.$

**THEOREM 2.** *Let  $T_n$  be a tournament with  $n \geq 3$  nodes such that  $u(T_n) = 3.$   
Then  $c(T_n) \leq C(n),$  where*

$$24C(n) = \begin{cases} n(n^2 - n + 2) & n = 6m \\ (n-1)(n^2 - 1) & n = 6m + 1 \\ n(n-2)(n+1) & n = 6m + 2 \\ n^3 - n^2 - n + 9 & \text{if } n = 6m + 3 \\ (n+2)(n^2 - 3n + 4) & n = 6m + 4 \\ (n-1)(n^2 - 1) & n = 6m + 5. \end{cases}$$

Furthermore,  $c(T_n) = C(n)$  if and only if  $T_n$  is a tournament of the form  
 $Q(R_X, R_Y, R_Z)$  where (i)  $R_X, R_Y,$  and  $R_Z$  are regular tournaments with  
 $X, Y,$  and  $Z$  nodes, (ii)  $X + Y + Z = n - 3,$  and (iii)  $X, Y,$  and  $Z$   
differ from each other by at most one.

**PROOF:** The theorem certainly holds when  $n = 3$  or  $4,$  so we may assume that  
 $n \geq 5.$  Let  $x, y,$  and  $z$  denote the uncovered nodes of the tournament  $T_n.$   
We may suppose that  $x \rightarrow y$  and  $y \rightarrow z.$  If  $x \rightarrow z$  then, since  $x$  does not  
cover  $z,$  there must be a fourth node  $v$  such that  $z \rightarrow v$  and  $v \rightarrow x;$  but  
then none of the nodes  $x, y,$  or  $z$  would cover  $v$  which, since the covering  
relation is transitive, contradicts the assumption that  $u(T_n) = 3.$  Consequently,  
 $z \rightarrow x$  and the uncovered nodes  $x, y,$  and  $z$  form a 3-cycle,  $(x, y, z)$  say.

Each node  $v \notin \{x, y, z\}$  is covered by at least one of the three uncovered  
nodes  $x, y,$  and  $z.$  If node  $x,$  say, covers such a node  $v$  then  $x \rightarrow v$  and,  
in addition,  $z \rightarrow v;$  for,  $z \rightarrow x$  and if  $v \rightarrow z,$  then  $x$  would not cover  $v.$   
Thus each such node  $v$  loses to at least two of the nodes  $x, y,$  and  $z,$  namely,  
a node that covers  $v$  and the immediate predecessor of the covering node in the  
3-cycle  $(x, y, z).$

Let  $T_X$  denote the (possibly empty) subtournament of  $T_n$  determined  
by those nodes  $v$  such that (i)  $v$  is covered by  $x$  and hence loses both  
to  $x$  and to  $z,$  the predecessor of  $x$  in the 3-cycle  $(x, y, z),$  but (ii)  $v$   
beats  $y,$  the successor of  $x$  in the 3-cycle  $(x, y, z).$  Let  $T_Y$  and  $T_Z$  be

similarly defined with respect to nodes  $y$  and  $z$  and, finally, let  $T_D$  denote the subtournament determined by the remaining nodes  $v$  that lose to all three nodes  $x, y$ , and  $z$ . Then  $X+Y+Z+D = n-3$  since each node  $v \notin \{x, y, z\}$  belongs to exactly one of these four subtournaments.

If there were nodes  $v \in T_X$  and  $w \in T_Y$ , say, such that  $w \rightarrow v$ , then  $(w, v, y)$  would be a 3-cycle containing the arc  $\overrightarrow{yw}$ , contrary to the assumption that  $y$  covers  $w$ . Consequently,  $T_X \rightarrow T_Y$  and, similarly,  $T_Y \rightarrow T_Z$  and  $T_Z \rightarrow T_X$ . The foregoing observations imply that  $T_n$  contains a subtournament  $Q(T_X, T_Y, T_Z)$  where  $X+Y+Z = n-3-D$  plus the (disjoint) subtournament  $T_D$  of nodes that lose to all three of the nodes  $x, y$ , and  $z$ .

Let  $s_x, s_y$ , and  $s_z$  denote the scores of the nodes  $x, y$ , and  $z$  in the tournament  $T_n$  and let  $s_1, s_2, \dots, s_{n-3}$  denote the scores of the remaining nodes. It follows from what we have deduced about the structure of  $T_n$  that

$$(6.1) \quad \begin{aligned} s_x + s_y + s_z &= (n-2-Z) + (n-2-X) + (n-2-Y) \\ &= 3(n-2) - (X+Y+Z) = 2n-3+D; \end{aligned}$$

furthermore,

$$(6.2) \quad s_1 + \dots + s_{n-3} = \binom{n}{2} - s_x - s_y - s_z = \binom{n-2}{2} - D.$$

It is not difficult to verify that the sequence  $\bar{s} = (s_1, \dots, s_{n-3}, s_x, s_y, s_z)$  satisfies the hypothesis of Lemma 2 with  $J = (n-2)(n-3)/2$ ,  $K = 2n-3$ ,  $j = n-3$ , and  $k = 3$ . Hence we conclude that a lower bound for the sum

$$f(\bar{s}) = \sum_1^{n-3} \binom{s_i}{2} + \binom{s_x}{2} + \binom{s_y}{2} + \binom{s_z}{2}$$

is obtained by evaluating the right hand side when the integers  $s_x, s_y, s_z$  are as nearly equal as possible and the integers  $s_1, \dots, s_{n-3}$  are as nearly equal as possible, subject to conditions (6.1) and (6.2) with  $D = 0$ .

More specifically, suppose that  $n = 6m + 5$ ; then  $2n - 3 = 12m + 7$  and

$$\binom{s_x}{2} + \binom{s_y}{2} + \binom{s_z}{2} \geq 2 \binom{4m+2}{2} + \binom{4m+3}{2} = (n-2)(2n-5)/3,$$

with equality holding if and only if  $s_x, s_y$ , and  $s_z$  equal  $4m+2, 4m+2$ , and  $4m+3$  or, equivalently,  $X, Y$ , and  $Z$  equal  $2m, 2m+1$ , and  $2m+1$  (in

some order). Furthermore,  $n - 2 = 6m + 3$  and  $n - 3 = 6m + 2$ , so

$$\sum_1^{n-3} \binom{s_i}{2} \geq (3m+1) \binom{3m+1}{2} + (3m+1) \binom{3m+2}{2} = (n-3)^3/8,$$

with equality holding if and only if half the scores  $s_1, \dots, s_{n-3}$  equal  $3m+1$  and the other half equal  $3m+2$ . Thus it follows from (2.1) that if  $n = 6m+5$ , then

$$\begin{aligned} c(T_n) &= \binom{n}{3} - f(\bar{s}) \\ &\leq \binom{n}{3} - (n-2)(2n-5)/3 - (n-3)^3/8 \\ &= (n^2-1)(n-1)/24 = C(6m+5). \end{aligned}$$

Moreover, equality holds if and only if  $D = 0$  and  $T_n$  is a tournament of the form  $Q(T_X, T_Y, T_Z)$  where  $X, Y$ , and  $Z$  equal  $2m, 2m+1$ , and  $2m+1$ ; and half the nodes in the subtournaments  $T_X, T_Y$ , and  $T_Z$  have score  $3m+1$  in  $T_n$  and the other half have score  $3m+2$ . It is not difficult to see that this last condition on the scores is satisfied if and only if all the subtournaments  $T_X, T_Y$ , and  $T_Z$  are regular. This suffices to prove the required result when  $n = 6m+5$ , and the same type of argument covers the cases  $n \equiv 0, 1$ , or  $3 \pmod{6}$  as well. (The cases  $n \equiv 0$  or  $3 \pmod{6}$ , are particularly easy.)

If, however,  $n \equiv 2$  or  $4 \pmod{6}$ , then the foregoing argument yields an upper bound for  $c(T_n)$  that is not best possible. For, to realize the bound in these cases, the nodes in the subtournaments  $T_X, T_Y$ , and  $T_Z$  would all have to have the same score and this is not possible here. Thus we need some additional arguments in these two remaining cases.

Suppose that  $n = 6m+2$  where  $m \geq 1$ . If  $D \geq 1$  for the tournament  $T_n$ , then it follows readily from Lemma 2 that

$$\begin{aligned} f(\bar{s}) &\geq \binom{4m}{2} + 2 \binom{4m+1}{2} + \binom{3m-1}{2} + (6m-2) \binom{3m}{2} \\ &= n(n-2)(3n-5)/24 + 1, \end{aligned}$$

so

$$c(T_n) = \binom{n}{3} - f(\bar{s}) \leq n(n-2)(n+1)/24 - 1 = C(6m+2) - 1.$$

Similarly, if  $n = 6m + 4$  and  $D \geq 1$ , we find that  $c(T_n) \leq C(6m + 4) - 1$ . So we may assume henceforth that  $n = 6m + 2$  or  $6m + 4$ , where  $m \geq 1$ , and that  $D = 0$ , that is, that  $T_n$  is a tournament of the form  $Q(T_X, T_Y, T_Z)$  where  $X + Y + Z = n - 3$ . We next dispose of the possibility that  $\max \{X, Y, Z\}$  exceeds  $4m$ , say.

If  $n$  is even - in particular, if  $n \equiv 2$  or  $4 \pmod{6}$  - then

$$\sum_1^{n-3} \binom{s_i}{2} \geq (n-3) \binom{(n-2)/2}{2} = (n-2)(n-3)(n-4)/8.$$

Now suppose that  $n = 6m + 2$ , where  $m \geq 1$ , so that  $X + Y + Z = 6m - 1$  and  $s_x + s_y + s_z = 2n - 3 = 12m + 1$ . If  $\max \{X, Y, Z\} \geq 4m + 1$ , then  $\min \{s_x, s_y, s_z\} \leq 2m - 1$  and

$$\binom{s_x}{2} + \binom{s_y}{2} + \binom{s_z}{2} \geq \binom{2m-1}{2} + 2 \binom{5m+1}{2} = (9n^2 - 32n + 40)/12,$$

appealing to Lemma 2 again. Hence, in this case,

$$\begin{aligned} c(T_n) &= \binom{n}{3} - f(\bar{s}) \\ &\leq \binom{n}{3} - (n-2)(n-3)(n-4)/8 - (9n^2 - 32n + 40)/12 \\ &= (n^3 - 3n^2 - 6n - 8)/24 \\ &= C(6m + 2) - (n^2 + 2n + 4)/12 < C(6m + 2). \end{aligned}$$

Similarly, we find that if  $n = 6m + 4$ , where  $m \geq 1$ , and  $\max \{X, Y, Z\} \geq 4m + 1$ , then

$$c(T_n) \leq (n^3 - 3n^2 + 10n - 8)/24 = C(6m + 4) - (n-2)(n-4)/12 < C(6m + 4).$$

Thus we may further assume, from now on, that  $\max \{X, Y, Z\} \leq 4m \leq 2(n-2)/3$ .

It follows readily from the definition of  $Q(T_X, T_Y, T_Z)$  that

$$c(T_n) = (X + 1)(Y + 1)(Z + 1) + c(T_X) + c(T_Y) + c(T_Z).$$

We mentioned earlier, in (2.2), that  $c(T_N) \leq \gamma(N)$ , where  $24\gamma(N) = N^3 - N$  or  $N^3 - 4N$  according as  $N$  is odd or even, with equality holding if and only

if  $T_N$  is a regular tournament  $R_N$ . Consequently,

$$c(T_n) \leq \Gamma(X, Y, Z)$$

where

$$\Gamma(X, Y, Z) = (X + 1)(Y + 1)(Z + 1) + \gamma(X) + \gamma(Y) + \gamma(Z),$$

with equality holding if and only if  $T_n$  is of the form  $Q(R_X, R_Y, R_Z)$ . So it remains to determine the values of  $X, Y$ , and  $Z$ , where  $X + Y + Z = n - 3$ , for which the function  $\Gamma(X, Y, Z)$  attains its maximum value. We need consider only the cases when  $Z \leq Y \leq X \leq 2(n - 2)/3$ .

Notice that it follows from the definition of the function  $\gamma(N)$  that

$$(6.3) \quad ((N - 1)^2 - 1)/8 \leq \gamma(N) - \gamma(N - 1) \leq (N^2 - 1)/8$$

for  $N = 1, 2, \dots$ , with equality holding on the left or the right according as  $N$  is even or odd. Now suppose that  $X > Z + 1$ . Then

$$\begin{aligned} \Delta &:= \Gamma(X - 1, Y, Z + 1) - \Gamma(X, Y, Z) \\ &= (X - Z - 1)(Y + 1) + \gamma(X - 1) - \gamma(X) + \gamma(Z + 1) - \gamma(Z) \\ &\geq (X - Z - 1)(Y + 1) - (X^2 - Z^2)/8 \\ &= (X - Z - 1)\{(Y + 1) - (X + Z + 1)/8\} - Z/4 - 1/8, \end{aligned}$$

where we have used relation (6.3) in the third line.

If  $X = Z + 2$ , then

$$\begin{aligned} \Delta &\geq Y + 1 - (2Z + 3)/8 - Z/4 - 1/8 \\ &= Y - Z/2 + 1/2 \geq Z/2 + 1/2 > 0. \end{aligned}$$

Next we combine the inequalities  $Z \leq Y$ ,  $Y \geq (n - 3 - X)/2$ , and  $X \leq 2(n - 2)/3$ , and find that

$$\begin{aligned} Y + 1 - (X + Z + 1)/8 &\geq (7Y + 7 - X)/8 \\ &\geq \{7(n - 1 - X) - 2X\}/16 \\ &\geq \{7(n - 1) - 6(n - 2)\}/16 = (n + 5)/16. \end{aligned}$$

Consequently, if  $X \geq Z + 3$ , then

$$\begin{aligned}\Delta &\geq (n + 5)/8 - Z/4 - 1/8 \\ &\geq (n + 4)/8 - (n - 3)/12 = n/24 + 3/4 > 0.\end{aligned}$$

Thus, if  $X > Z + 1$ , then  $\Gamma(X - 1, Y, Z + 1) > \Gamma(X, Y, Z)$ . This implies that if  $n \equiv 2$  or  $4 \pmod{6}$  and  $\max\{X, Y, Z\} \leq 2(n - 2)/3$ , then the maximum value of  $\Gamma(X, Y, Z)$  occurs when  $X, Y$ , and  $Z$  are as nearly equal as possible. It is easy to verify that this maximum value of  $\Gamma(X, Y, Z)$  equals  $C(n)$  when  $n \equiv 2$  or  $4 \pmod{6}$  (and, in fact, for all  $n$ ), so this suffices to complete the proof of the theorem.

Let  $M_n$  and  $Q_n$  denote any of the minimal and maximal tournaments considered in Theorems 1 and 2, respectively. We remark in closing that if  $n = 3$  or  $n \geq 5$ , then there exists a tournament  $T_n$  such that  $c(T_n) = c(M_n)$  but for which  $u(T_n) = 3$  whereas  $u(M_n) = n$ . Furthermore, if  $n \geq 18$  (and perhaps for some smaller values also), then there exist tournaments  $T_n$  such that  $c(T_n) = c(Q_n)$  but  $u(T_n) = n - 3$  whereas  $u(Q_n) = 3$ .

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